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QSPEC Template for the Quality-of-Service NSIS Signaling Layer Protocol (NSLP)

Abstract

The Quality-of-Service (QoS) NSIS signaling layer protocol (NSLP) is used to signal QoS reservations and is independent of a specific QoS model (QOSM) such as IntServ or Diffserv. Rather, all information specific to a QOSM is encapsulated in a separate object, the QSPEC. This document defines a template for the QSPEC including a number of QSPEC parameters. The QSPEC parameters provide a common language to be reused in several QOSMs and thereby aim to ensure the extensibility and interoperability of QoS NSLP. While the base protocol is QOSM-agnostic, the parameters that can be carried in the QSPEC object are possibly closely coupled to specific models. The node initiating the NSIS signaling adds an Initiator QSPEC, which indicates the QSPEC parameters that must be interpreted by the downstream nodes less the reservation fails, thereby ensuring the intention of the NSIS initiator is preserved along the signaling path.

Status of This Memo

This document is not an Internet Standards Track specification; it is published for examination, experimental implementation, and evaluation.

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Ash, et al.

Experimental

[Page 1]

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Ash, et al. Experimental

[Page 2]

Table of Contents

1.	Introduction
	1.1. Conventions Used in This Document
2.	Terminology6
3.	QSPEC Framework7
	3.1. QoS Models
	3.2. QSPEC Objects
	3.3. QSPEC Parameters11
	3.3.1. Traffic Model Parameter12
	3.3.2. Constraints Parameters14
	3.3.3. Traffic-Handling Directives
	3.3.4. Traffic Classifiers17
	3.4. Example of QSPEC Processing17
4.	QSPEC Processing and Procedures
	4.1. Local QSPEC Definition and Processing20
	4.2. Reservation Success/Failure, QSPEC Error Codes,
	and INFO-SPEC Notification23
	4.2.1. Reservation Failure and Error E Flag
	4.2.2. QSPEC Parameter Not Supported N Flag25
	4.2.3. INFO-SPEC Coding of Reservation Outcome
	4.2.4. QNE Generation of a RESPONSE Message
	4.2.5. Special Case of Local QSPEC
	4.3. QSPEC Procedures27
	4.3.1. Two-Way Transactions
	4.3.2. Three-Way Transactions
	4.3.3. Resource Queries
	4.3.4. Bidirectional Reservations
	4.3.5. Preemption
_	4.4. QSPEC Extensibility
5.	QSPEC Functional Specification
	5.1. General QSPEC Formats
	5.1.1. Common Header Format
	5.1.2. QSPEC Object Header Format
	5.2. QSPEC Parameter Coding
	5.2.1. <tmod-1> Parameter</tmod-1>
	5.2.2. <tmod-2> Parameter</tmod-2>
	5.2.3. <path latency=""> Parameter</path>
	5.2.4. <path jitter=""> Parameter40 5.2.5. <path plr=""> Parameter41</path></path>
	5.2.5. <path per=""> Parameter</path>
	5.2.6. <path per=""> Parameter</path>
	5.2.7. <stack ferms="" parameter43<br="">5.2.8. <preemption priority=""> and <defending priority=""></defending></preemption></stack>
	5.2.8. <preemption priority=""> and <detending priority=""> Parameters</detending></preemption>
	5.2.9. <admission priority=""> Parameter43</admission>
	5.2.9. <addission priority=""> Parameter44 5.2.10. <rph priority=""> Parameter45</rph></addission>
	5.2.10. <rph parameter45<br="" priority'="">5.2.11. <excess treatment=""> Parameter46</excess></rph>
	5.2.12. <phb class=""> Parameter</phb>
	5.2.12. SPAB CLASS? PALAMELET48

Ash, et al.

Experimental

[Page 3]

5.2.13. <dste class="" type=""> Parameter</dste>
5.2.14. <y.1541 class="" qos=""> Parameter</y.1541>
6. Security Considerations51
7. IANA Considerations
8. Acknowledgements
9. Contributors
10. Normative References57
11. Informative References59
Appendix A. Mapping of QoS Desired, QoS Available, and QoS
Reserved of NSIS onto AdSpec, TSpec, and RSpec of RSVP IntServ .62
Appendix B. Example of TMOD Parameter Encoding62

1. Introduction

The QoS NSIS signaling layer protocol (NSLP) [RFC5974] is used to signal QoS reservations for a data flow, provide forwarding resources (QoS) for that flow, and establish and maintain state at nodes along the path of the flow. The design of QoS NSLP is conceptually similar to the decoupling between RSVP [RFC2205] and the IntServ architecture [RFC2210], where a distinction is made between the operation of the signaling protocol and the information required for the operation of the Resource Management Function (RMF). [RFC5974] describes the signaling protocol, while this document describes the RMF-related information carried in the QSPEC (QoS Specification) object carried in QoS NSLP messages.

[RFC5974] defines four QoS NSLP messages -- RESERVE, QUERY, RESPONSE, and NOTIFY -- each of which may carry the QSPEC object, while this document describes a template for the QSPEC object. The QSPEC object carries information on traffic descriptions, resources required, resources available, and other information required by the RMF. Therefore, the QSPEC template described in this document is closely tied to QoS NSLP, and the reader should be familiar with [RFC5974] to fully understand this document.

A QoS-enabled domain supports a particular QoS model (QOSM), which is a method to achieve QoS for a traffic flow. A QOSM incorporates QoS provisioning methods and a QoS architecture, and defines the behavior of the RMF that reserves resources for each flow, including inputs and outputs. The QoS NSLP protocol is able to signal QoS reservations for different QOSMs, wherein all information specific to a QOSM is encapsulated in the QSPEC object, and only the RMF specific to a given QOSM will need to interpret the QSPEC. Examples of QOSMs are IntServ, Diffserv admission control, and those specified in [CL-QOSM], [RFC5976], and [RFC5977].

Ash, et al. Experimental

[Page 4]

QSPEC parameters include, for example:

- o a mandatory traffic model (TMOD) parameter,
- o constraints parameters such as path latency and path jitter,
- o traffic handling directives such as excess treatment, and
- o traffic classifiers such as PHB class.

While the base protocol is QOSM-agnostic, the parameters that can be carried in the QSPEC object are possibly closely coupled to specific models.

QSPEC objects loosely correspond to the TSpec, RSpec, and AdSpec objects specified in RSVP and may contain, respectively, a description of QoS Desired, QoS Reserved, and QoS Available. Going beyond RSVP functionality, the QSPEC also allows indicating a range of acceptable QoS by defining a QSPEC object denoting minimum QoS. Usage of these QSPEC objects is not bound to particular message types, thus allowing for flexibility. A QSPEC object collecting information about available resources may travel in any QoS NSLP message, for example, a QUERY message or a RESERVE message, as defined in [RFC5974]. The QSPEC travels in QoS NSLP messages but is opaque to the QoS NSLP and is only interpreted by the RMF.

Interoperability between QoS NSIS entities (QNEs) in different domains is enhanced by the definition of a common set of QSPEC parameters. A QoS NSIS initiator (QNI) initiating the QoS NSLP signaling adds an Initiator QSPEC object containing parameters describing the desired QoS, normally based on the QOSM it supports. QSPEC parameters flagged by the QNI must be interpreted by all QNEs in the path, else the reservation fails. In contrast, QSPEC parameters not flagged by the QNI may be skipped if not understood. Additional QSPEC parameters can be defined by informational specification documents, and thereby ensure the extensibility and flexibility of QoS NSLP.

A Local QSPEC can be defined in a local domain with the Initiator QSPEC encapsulated, where the Local QSPEC must be functionally consistent with the Initiator QSPEC in terms of defined source traffic and other constraints. That is, a domain-specific local QSPEC can be defined and processed in a local domain, which could, for example, enable simpler processing by QNEs within the local domain.

In Section 3.4, an example of QSPEC processing is provided.

Ash, et al. Experimental

[Page 5]

RFC 5975

1.1. Conventions Used in This Document

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119 [RFC2119].

2. Terminology

Initiator QSPEC: The Initiator QSPEC is included in a QoS NSLP message by the QNI/QNR. It travels end-to-end to the QNR/QNI and is never removed.

Local QSPEC: A Local QSPEC is used in a local domain and is domain specific. It encapsulates the Initiator QSPEC and is removed at the egress of the local domain.

Minimum QoS: QSPEC object that, together with a description of QoS Desired or QoS Available, allows the QNI to specify a QoS range, i.e., an upper and lower bound. If the QoS Desired cannot be reserved, QNEs are going to decrease the reservation until the minimum QoS is hit. Note that the term "minimum" is used generically, since for some parameters, such as loss rate and latency, what is specified is the maximum acceptable value.

QNE: QOS NSIS Entity, a node supporting QoS NSLP.

QNI: QOS NSIS Initiator, a node initiating QoS NSLP signaling.

QNR: QOS NSIS Receiver, a node terminating QoS NSLP signaling.

QoS Available: QSPEC object containing parameters describing the available resources. They are used to collect information along a reservation path.

QoS Desired: QSPEC object containing parameters describing the desired QoS for which the sender requests reservation.

QoS Model (QOSM): a method to achieve QoS for a traffic flow, e.g., IntServ Controlled Load; specifies the subset of QSPEC QoS constraints and traffic handling directives that a QNE implementing that QOSM is capable of supporting and how resources will be managed by the RMF.

QoS Reserved: QSPEC object containing parameters describing the reserved resources and related QoS parameters.

QSPEC: the object of QoS NSLP that contains all QoS-specific information.

Ash, et al. Experimental

[Page 6]

QSPEC parameter: Any parameter appearing in a QSPEC; for example, traffic model (TMOD), path latency, and excess treatment parameters.

QSPEC Object: Main building blocks containing a QSPEC parameter set that is the input or output of an RMF operation.

QSPEC Type: Identifies a particular QOSM used in the QSPEC

Resource Management Function (RMF): Functions that are related to resource management and processing of QSPEC parameters.

3. QSPEC Framework

The overall framework for the QoS NSLP is that [RFC5974] defines QoS signaling and semantics, the QSPEC template defines the container and semantics for QoS parameters and objects, and informational specifications define QoS methods and procedures for using QoS signaling and QSPEC parameters/objects within specific QoS deployments. QoS NSLP is a generic QoS signaling protocol that can signal for many QOSMs.

3.1. QoS Models

A QOSM is a method to achieve QoS for a traffic flow, e.g., IntServ Controlled Load [CL-QOSM], Resource Management with Diffserv [RFC5977], and QoS signaling for Y.1541 QoS classes [RFC5976]. A QOSM specifies a set of QSPEC parameters that describe the QoS desired and how resources will be managed by the RMF. The RMF implements functions that are related to resource management and processes the QSPEC parameters.

QOSMs affect the operation of the RMF in NSIS-capable nodes and the information carried in QSPEC objects. Under some circumstances (e.g., aggregation), they may cause a separate NSLP session to be instantiated by having the RMF as a QNI. QOSM specifications may define RMF triggers that cause the QoS NSLP to run semantics within the underlying QoS NSLP signaling state and messaging processing rules, as defined in Section 5.2 of [RFC5974]. New QoS NSLP message processing rules can only be defined in extensions to QoS NSLP. If a QOSM specification defines triggers that deviate from existing QoS NSLP processing rules, the fallback for QNEs not supporting that QOSM are the QoS NSLP state transition/message processing rules.

The QOSM specification includes how the requested QoS resources will be described and how they will be managed by the RMF. For this purpose, the QOSM specification defines a set of QSPEC parameters it uses to describe the desired QoS and resource control in the RMF, and it may define additional QSPEC parameters.

Ash, et al. Experimental

[Page 7]

When a QoS NSLP message travels through different domains, it may encounter different QOSMs. Since QOSMs use different QSPEC parameters for describing resources, the QSPEC parameters included by the QNI may not be understood in other domains. The QNI therefore can flag those QSPEC parameters it considers vital with the M flag. QSPEC parameters with the M flag set must be interpreted by the downstream QNEs, or the reservation fails. QSPEC parameters without the M flag set should be interpreted by the downstream QNEs, but may be ignored if not understood.

A QOSM specification SHOULD include the following:

- role of QNEs, e.g., location, frequency, statefulness, etc.
- QSPEC definition including QSPEC parameters
- QSPEC procedures applicable to this QOSM
- QNE processing rules describing how QSPEC information is treated and interpreted in the RMF, e.g., admission control, scheduling, policy control, QoS parameter accumulation (e.g., delay)
- at least one bit-level QSPEC example
- QSPEC parameter behavior for new QSPEC parameters that the QOSM specification defines
- a definition of what happens in case of preemption if the default QNI behavior (teardown preempted reservation) is not followed (see Section 4.3.5)
- A QOSM specification MAY include the following:
- definitions of additional QOSM-specific error codes, as discussed in Section 4.2.3
- the QoS-NSLP options a QOSM wants to use, when several options are available for a QOSM (e.g., Local QSPEC to either a) hide the Initiator QSPEC within a local domain message, or b) encapsulate the Initiator QSPEC).

QOSMs are free, subject to IANA registration and review rules, to extend QSPECs by adding parameters of any of the kinds supported by the QSPEC. This includes traffic description parameters, constraint parameters, and traffic handling directives. QOSMs are not permitted, however, to reinterpret or redefine the QSPEC parameters specified in this document. Note that signaling functionality is only defined by the QoS NSLP document [RFC5974] and not by this document or by QOSM specification documents.

Ash, et al. Experimental

[Page 8]

3.2. QSPEC Objects

The QSPEC is the object of QoS NSLP containing QSPEC objects and parameters. QSPEC objects are the main building blocks of the QSPEC parameter set that is input or output of an RMF operation. QSPEC parameters are the parameters appearing in a QSPEC, which must include the traffic model parameter (TMOD), and may optionally include constraints (e.g., path latency), traffic handling directives (e.g., excess treatment), and traffic classifiers (e.g., PHB class). The RMF implements functions that are related to resource management and processes the QSPEC parameters.

The QSPEC consists of a QSPEC version number and QSPEC objects. IANA assigns a new QSPEC version number when the current version is deprecated or deleted (as required by a specification). Note that a new QSPEC version number is not needed when new QSPEC parameters are specified. Later QSPEC versions MUST be backward compatible with earlier QSPEC versions. That is, a version n+1 device must support QSPEC version n (or earlier). On the other hand, if a QSPEC version n (or earlier) device receives an NSLP message specifying QSPEC version n+1, then the version n device responds with an 'Incompatible QSPEC' error code (0x0f) response, as discussed in Section 4.2.3, allowing the QNE that sent the NSLP message to retry with a lower QSPEC version.

This document provides a template for the QSPEC in order to promote interoperability between QOSMs. Figure 1 illustrates how the QSPEC is composed of up to 4 QSPEC objects, namely QoS Desired, QoS Available, QoS Reserved, and Minimum QoS. Each of these QSPEC objects consists of a number of QSPEC parameters. A given QSPEC may contain only a subset of the QSPEC objects, e.g., QoS Desired. The QSPEC objects QoS Desired, QoS Available, QoS Reserved and Minimum QoS MUST all be supported by QNEs and MAY appear in any QSPEC object carried in any QoS NSLP message (RESERVE, QUERY, RESPONSE, NOTIFY). See [RFC5974] for descriptions of the QoS NSLP RESERVE, QUERY, RESPONSE, and NOTIFY messages.

Ash, et al. Experimental

[Page 9]



Figure 1: Structure of the QSPEC

Use of the 4 QSPEC objects (QoS Desired, QoS Available, QoS Reserved, and Minimum QoS) is described in Section 4.3 for 3 message sequences and 7 object combinations.

The QoS Desired Object describe the resources the QNI desires to reserve, and hence this is a read-only QSPEC object in that the QSPEC parameters carried in the object may not be overwritten. QoS Desired is always included in a RESERVE message and sometimes included in the QUERY message (see Section 4.3 for details).

As described in Section 4.3, the QoS Available object may travel in a RESERVE message, RESPONSE Message, or QUERY message and may collect information on the resources currently available on the path. In this case, QoS Available is a read-write object, which means the QSPEC parameters contained in QoS Available may be updated, but they cannot be deleted. As such, each QNE MUST inspect all parameters of this QSPEC object, and if resources available to this QNE are less than what a particular parameter says currently, the QNE MUST adapt this parameter accordingly. Hence, when the message arrives at the recipient of the message, <QoS Available> reflects the bottleneck of the resources currently available on a path. It can be used in a QUERY message, for example, to collect the available resources along a data path.

When QoS Available travels in a RESPONSE message, it in fact just transports the result of a previous measurement performed by a RESERVE or QUERY message back to the initiator. Therefore, in this

Ash, et al.

Experimental

[Page 10]

case, QoS Available is read-only. In one other instance described in Section 4.3.2 (Case 3), QoS Available is sent by the QNI in a RESERVE message as a read-only QSPEC object (see Section 4.3.2 for details).

The QoS Reserved object reflects the resources that are being reserved. It is a read-only object and is always included in a RESPONSE message if QoS Desired is included in the RESERVE message (see Section 4.3 for details).

Minimum QoS does not have an equivalent in RSVP. It allows the QNI to define a range of acceptable QoS levels by including both the desired QoS value and the minimum acceptable QoS in the same message. Note that the term "minimum" is used generically, since for some parameters, such as loss rate and latency, what is specified is the maximum acceptable value. It is a read-only object, and may be included in a RESERVE message, RESPONSE message, or QUERY message (see Section 4.3 for details). The desired QoS is included with a QoS Desired and/or a QoS Available QSPEC object seeded to the desired QoS value. The minimum acceptable QoS value MAY be coded in the Minimum QoS QSPEC object. As the message travels towards the QNR, QoS Available is updated by QNEs on the path. If its value drops below the value of Minimum QoS, the reservation fails and is aborted. When this method is employed, the QNR signals back to the QNI the value of QoS Available attained in the end, because the reservation may need to be adapted accordingly (see Section 4.3 for details).

Note that the relationship of QSPEC objects to RSVP objects is covered in Appendix A.

3.3. QSPEC Parameters

QSPEC parameters provide a common language for building QSPEC objects. This document defines a number of QSPEC parameters; additional parameters may be defined in separate QOSM specification documents. For example, QSPEC parameters are defined in [RFC5976] and [RFC5977].

One QSPEC parameter, <TMOD>, is special. It provides a description of the traffic for which resources are reserved. This parameter must be included by the QNI, and it must be interpreted by all QNEs. All other QSPEC parameters are populated by a QNI if they are applicable to the underlying QoS desired. For these QSPEC parameters, the QNI sets the M flag if they must be interpreted by downstream QNEs. If QNEs cannot interpret the parameter, the reservation fails. QSPEC parameters populated by a QNI without the M flag set should be interpreted by downstream QNEs, but may be ignored if not understood.

Ash, et al. Experimental

[Page 11]

In this document, the term 'interpret' means, in relation to RMF processing of QSPEC parameters, that the RMF processes the QSPEC parameter according to the commonly accepted normative procedures specified by references given for each QSPEC parameter. Note that a QNE need only interpret a QSPEC parameter if it is populated in the QSPEC object by the QNI; if not populated in the QSPEC, the QNE does not interpret it of course.

Note that when an ingress QNE in a local domain defines a Local QSPEC and encapsulates the Initiator QSPEC, the QNEs in the interior local domain need only process the Local QSPEC and can ignore the Initiator (encapsulated) QSPEC. However, edge QNEs in the local domain indeed must interpret the QSPEC parameters populated in the Initiator QSPEC with the M flag set and should interpret QSPEC parameters populated in the Initiator QSPEC without the M flag set.

As described in the previous section, QoS parameters may be overwritten depending on which QSPEC object and which message they appear in.

3.3.1. Traffic Model Parameter

The <Traffic Model> (TMOD) parameter is mandatory for the QNI to include in the Initiator QSPEC and mandatory for downstream QNEs to interpret. The traffic description specified by the TMOD parameter is a container consisting of 5 sub-parameters [RFC2212]:

- o rate (r) specified in octets per second
- o bucket size (b) specified in octets
- o peak rate (p) specified in octets per second
- o minimum policed unit (m) specified in octets
- o maximum packet size (MPS) specified in octets

The TMOD parameter takes the form of a token bucket of rate (r) and bucket size (b), plus a peak rate (p), minimum policed unit (m), and maximum packet size (MPS).

Both b and r MUST be positive. The rate, r, is measured in octets of IP packets per second, and can range from 1 octet per second to as large as 40 teraoctets per second. The bucket depth, b, is also measured in octets and can range from 1 octet to 250 gigaoctets. The peak rate, p, is measured in octets of IP packets per second and has the same range and suggested representation as the bucket rate.

The peak rate is the maximum rate at which the source and any reshaping (defined below) may inject bursts of traffic into the network. More precisely, it is a requirement that for all time periods the amount of data sent cannot exceed MPS+pT, where MPS is

Ash, et al. Experimental

[Page 12]

the maximum packet size and T is the length of the time period. Furthermore, p MUST be greater than or equal to the token bucket rate, r. If the peak rate is unknown or unspecified, then p MUST be set to infinity.

The minimum policed unit, m, is an integer measured in octets. All IP packets less than size m will be counted, when policed and tested for conformance to the TMOD, as being of size m.

The maximum packet size, MPS, is the biggest packet that will conform to the traffic specification; it is also measured in octets. The flow MUST be rejected if the requested maximum packet size is larger than the MTU of the link. Both m and MPS MUST be positive, and m MUST be less than or equal to MPS.

Policing compares arriving traffic against the TMOD parameters at the edge of the network. Traffic is policed to ensure it conforms to the token bucket. Reshaping attempts to restore the (possibly distorted) traffic's shape to conform to the TMOD parameters, and traffic that is in violation of the TMOD is discovered because the reshaping fails and the reshaping buffer overflows.

The token bucket and peak rate parameters require that traffic MUST obey the rule that over all time periods, the amount of data sent cannot exceed MPS+min[pT, rT+b-MPS], where r and b are the token bucket parameters, MPS is the maximum packet size, and T is the length of the time period (note that when p is infinite, this reduces to the standard token bucket requirement). For the purposes of this accounting, links MUST count packets that are smaller than the minimum policing unit as being of size m. Packets that arrive at an element and cause a violation of the MPS + min[pT, rT+b-MPS] bound are considered non-conformant.

All 5 of the sub-parameters MUST be included in the TMOD parameter. The TMOD parameter can be set to describe the traffic source. If, for example, TMOD is set to specify bandwidth only, then set r = peakrate = p, b = large, and m = large. As another example, if TMOD is set for TCP traffic, then set r = average rate, b = large, and p =large.

When the 5 TMOD sub-parameters are included in QoS Available, they provide information, for example, about the TMOD resources available along the path followed by a data flow. The value of TMOD at a QNE is an estimate of the TMOD resources the QNE has available for packets following the path up to the next QNE, including its outgoing link, if this link exists. Furthermore, the QNI MUST account for the resources of the ingress link, if this link exists. Computation of

Ash, et al. Experimental

[Page 13]

the value of this parameter SHOULD take into account all information available to the QNE about the path, taking into consideration administrative and policy controls, as well as physical resources.

The output composed value is the minimum of the QNE's value and the input composed value for r, b, p, and MPS, and the maximum of the QNE's value and the input composed value for m. This quantity, when composed end-to-end, informs the QNR (or QNI in a RESPONSE message) of the minimal TMOD resources along the path from QNI to QNR.

Two TMOD parameters are defined in Section 5, <TMOD-1> and <TMOD-2>, where the second parameter (<TMOD-2>) is specified as could be needed to support some Diffserv applications. For example, it is typically assumed that Diffserv Expedited Forwarding (EF) traffic is shaped at the ingress by a single rate token bucket. Therefore, a single TMOD parameter is sufficient to signal Diffserv EF traffic. However, for Diffserv Assured Forwarding (AF) traffic, two sets of token bucket parameters are needed -- one for the average traffic and one for the burst traffic. [RFC2697] defines a Single Rate Three Color Marker (srTCM), which meters a traffic stream and marks its packets according to three traffic parameters, Committed Information Rate (CIR), Committed Burst Size (CBS), and Excess Burst Size (EBS), to be either green, yellow, or red. A packet is marked green if it does not exceed the CBS; yellow if it does exceed the CBS, but not the EBS; and red otherwise. [RFC2697] defines specific procedures using two token buckets that run at the same rate. Therefore, 2 TMOD parameters are sufficient to distinguish among 3 levels of drop precedence. An example is also described in the Appendix to [RFC2597].

3.3.2. Constraints Parameters

<Path Latency>, <Path Jitter>, <Path PLR>, and <Path PER> are QSPEC parameters describing the desired path latency, path jitter, packet loss ratio, and path packet error ratio, respectively. Since these parameters are cumulative, an individual QNE cannot decide whether the desired path latency, etc., is available, and hence they cannot decide whether a reservation fails. Rather, when these parameters are included in <Desired QoS>, the QNI SHOULD also include corresponding parameters in a QoS Available QSPEC object in order to facilitate collecting this information.

The <Path Latency> parameter accumulates the latency of the packet forwarding process associated with each QNE, where the latency is defined to be the mean packet delay, measured in microseconds, added by each QNE. This delay results from the combination of link propagation delay, packet processing, and queuing. Each QNE MUST add the propagation delay of its outgoing link, if this link exists.

Ash, et al. Experimental

[Page 14]

Furthermore, the QNI SHOULD add the propagation delay of the ingress link, if this link exists. The composition rule for the <Path Latency> parameter is summation with a clamp of (2^{32}) - 1 on the maximum value. This quantity, when composed end-to-end, informs the QNR (or QNI in a RESPONSE message) of the minimal packet delay along the path from QNI to QNR. The purpose of this parameter is to provide a minimum path latency for use with services that provide estimates or bounds on additional path delay [RFC2212].

The <Path Jitter> parameter accumulates the jitter of the packet forwarding process associated with each QNE, where the jitter is defined to be the nominal jitter, measured in microseconds, added by each QNE. IP packet jitter, or delay variation, is defined in [RFC3393], Section 3.4 (Type-P-One-way-ipdv), and where the [RFC3393] selection function includes the packet with minimum delay such that the distribution is equivalent to 2-point delay variation in [Y.1540]. The suggested evaluation interval is 1 minute. This jitter results from packet-processing limitations, and includes any variable queuing delay that may be present. Each QNE MUST add the jitter of its outgoing link, if this link exists. Furthermore, the QNI SHOULD add the jitter of the ingress link, if this link exists. The composition method for the <Path Jitter> parameter is the combination of several statistics describing the delay variation distribution with a clamp on the maximum value (note that the methods of accumulation and estimation of nominal QNE jitter are specified in clause 8 of [Y.1541]). This quantity, when composed end-to-end, informs the QNR (or QNI in a RESPONSE message) of the nominal packet jitter along the path from QNI to QNR. The purpose of this parameter is to provide a nominal path jitter for use with services that provide estimates or bounds on additional path delay [RFC2212].

The <Path PLR> parameter is the unit-less ratio of total lost IP packets to total transmitted IP packets. <Path PLR> accumulates the packet loss ratio (PLR) of the packet-forwarding process associated with each QNE, where the PLR is defined to be the PLR added by each QNE. Each QNE MUST add the PLR of its outgoing link, if this link exists. Furthermore, the QNI MUST add the PLR of the ingress link, if this link exists. The composition rule for the <Path PLR> parameter is summation with a clamp on the maximum value. (This assumes sufficiently low PLR values such that summation error is not significant; however, a more accurate composition function is specified in clause 8 of [Y.1541].) This quantity, when composed end-to-end, informs the QNR (or QNI in a RESPONSE message) of the minimal packet PLR along the path from QNI to QNR.

Packet error ratio [Y.1540, Y.1541] is the unit-less ratio of total errored IP packet outcomes to the total of successful IP packet transfer outcomes plus errored IP packet outcomes in a population of

Ash, et al. Experimental

[Page 15]

interest, with a resolution of at least 10⁻⁹. If lesser resolution is available in a value, the unused digits MUST be set to zero. Note that the number of errored packets observed is directly related to the confidence in the result. The <Path PER> parameter accumulates the packet error ratio (PER) of the packet forwarding process associated with each QNE, where the PER is defined to be the PER added by each QNE. Each QNE MUST add the PER of its outgoing link, if this link exists. Furthermore, the QNI SHOULD add the PER of the ingress link, if this link exists. The composition rule for the <Path PER> parameter is summation with a clamp on the maximum value. (This assumes sufficiently low PER values such that summation error is not significant; however, a more accurate composition function is specified in clause 8 of [Y.1541].) This quantity, when composed end-to-end, informs the QNR (or QNI in a RESPONSE message) of the minimal packet PER along the path from QNI to QNR.

The slack term parameter is the difference between desired delay and delay obtained by using bandwidth reservation, and it is used to reduce the resource reservation for a flow [RFC2212].

3.3.3. Traffic-Handling Directives

An application MAY like to reserve resources for packets and also specify a specific traffic-handling behavior, such as <Excess Treatment>. In addition, as discussed in Section 3.1, an application MAY like to define RMF triggers that cause the QoS NSLP to run semantics within the underlying QoS NSLP signaling state / messaging processing rules, as defined in Section 5.2 of [RFC5974]. Note, however, that new QoS NSLP message processing rules can only be defined in extensions to the QoS NSLP. As with constraints parameters and other QSPEC parameters, Traffic Handling Directives parameters may be defined in QOSM specifications in order to provide support for QOSM-specific resource management functions. Such QOSMspecific parameters are already defined, for example, in [RFC5976], [RFC5977], and [CL-QOSM]. Generally, a Traffic Handling Directives parameters is expected to be set by the QNI in <QoS Desired>, and to not be included in <QoS Available>. If such a parameter is included in <QoS Available>, QNEs may change their value.

The <Preemption Priority> parameter is the priority of the new flow compared with the <Defending Priority> of previously admitted flows. Once a flow is admitted, the preemption priority becomes irrelevant. The <Defending Priority> parameter is used to compare with the preemption priority of new flows. For any specific flow, its preemption priority MUST always be less than or equal to the defending priority. <Admission Priority> and <RPH Priority> provide an essential way to differentiate flows for emergency services, Emergency Telecommunications Service (ETS), E911, etc., and assign

Ash, et al. Experimental

[Page 16]

them a higher admission priority than normal priority flows and besteffort priority flows.

The <Excess Treatment> parameter describes how the QNE will process out-of-profile traffic. Excess traffic MAY be dropped, shaped, and/or re-marked.

3.3.4. Traffic Classifiers

An application MAY like to reserve resources for packets with a particular Diffserv per-hop behavior (PHB) [RFC2475]. Note that PHB class is normally set by a downstream QNE to tell the QNI how to mark traffic to ensure the treatment that is designated by admission control; however, setting of the parameter by the QNI is not precluded. An application MAY like to reserve resources for packets with a particular QoS class, e.g., Y.1541 QoS class [Y.1541] or Diffserv-aware MPLS traffic engineering (DSTE) class type [RFC3564, RFC4124]. These parameters are useful in various QOSMs, e.g., [RFC5976], [RFC5977], and other QOSMs yet to be defined (e.g., DSTE-QOSM). This is intended to provide guidelines to QOSMs on how to encode these parameters; use of the PHB class parameter is illustrated in the example in the following section.

3.4. Example of QSPEC Processing

This section illustrates the operation and use of the QSPEC within the NSLP. The example configuration in shown in Figure 2.



Figure 2: Example Configuration of QoS-NSLP/QSPEC Operation

In this configuration, a laptop computer and a handheld wireless device are the endpoints for some application that has QoS requirements. Assume initially that the two endpoints are stationary during the application session, later we consider mobile endpoints.

Ash, et al. Experimental

[Page 17]

For this session, the laptop computer is connected to a home network that has no QoS support. The home network is connected to a CableLabs-type cable access network with dynamic QoS (DQOS) support, such as specified in the [DQOS] for cable access networks. That network is connected to a Diffserv core network that uses the Resource Management in Diffserv QoS Model [RFC5977]. On the other side of the Diffserv core is a wireless access network built on generation "X" technology with QoS support as defined by generation "X". And finally, the handheld endpoint is connected to the wireless access network.

We assume that the laptop is the QNI, and the handheld device is the QNR. The QNI will signal an Initiator QSPEC object to achieve the QoS desired on the path.

The QNI sets QoS Desired, QoS Available, and possibly Minimum QoS QSPEC objects in the Initiator QSPEC, and initializes QoS Available to QoS Desired. Each QNE on the path reads and interprets those parameters in the Initiator QSPEC and checks to see if QoS Available resources can be reserved. If not, the QNE reduces the respective parameter values in QoS Available and reserves these values. The minimum parameter values are given in Minimum QoS, if populated; they are zero if Minimum QoS is not included. If one or more parameters in QoS Available fails to satisfy the corresponding minimum values in Minimum QoS, the QNE generates a RESPONSE message to the QNI and the reservation is aborted. Otherwise, the QNR generates a RESPONSE to the QNI with the QoS Available for the reservation. If a QNE cannot reserve QoS Desired resources, the reservation fails.

The QNI populates QSPEC parameters to ensure correct treatment of its traffic in domains down the path. Let us assume the QNI wants to achieve QoS guarantees similar to IntServ Controlled Load service, and also is interested in what path latency it can achieve. Additionally, to ensure correct treatment further down the path, the QNI includes <PHB Class> in <QoS Desired>. The QNI therefore includes in the QSPEC

QoS Desired = <TMOD> <PHB Class> QoS Available = <TMOD> <Path Latency>

Since <Path Latency> and <PHB Class> are not vital parameters from the QNI's perspective, it does not raise their M flags.

There are three possibilities when a RESERVE message is received at a QNE at a domain border; they are described in the example:

- the QNE just leaves the QSPEC as is.

Ash, et al. Experimental

[Page 18]

- the QNE can add a Local QSPEC and encapsulate the Initiator QSPEC (see discussion in Section 4.1; this is new in QoS NSLP -- RSVP does not do this).
- the QNE can 'hide' the initiator RESERVE message so that only the edge QNE processes the initiator RESERVE message, which then bypasses intermediate nodes between the edges of the domain and issues its own local RESERVE message (see Section 3.3.1 of [RFC5974]). For this new local RESERVE message, the QNE acts as the QNI, and the QSPEC in the domain is an Initiator QSPEC. A similar procedure is also used by RSVP in making aggregate reservations, in which case there is not a new intra-domain (aggregate) RESERVE for each newly arriving inter-domain (per-flow) RESERVE, but the aggregate reservation is updated by the border QNE (or QNI) as need be. This is also how RMD works [RFC5977].

For example, at the RMD domain, a local RESERVE with its own RMD Initiator QSPEC corresponding to the RMD-QOSM is generated based on the original Initiator QSPEC according to the procedures described in Section 4.5 of [RFC5974] and in [RFC5977]. The ingress QNE to the RMD domain maps the TMOD parameters contained in the original Initiator QSPEC to the equivalent TMOD parameter representing only the peak bandwidth in the Local QSPEC. The local RMD QSPEC for example also needs <PHB Class>, which in this case was provided by the QNI.

Furthermore, if the node can, at the egress to the RMD domain, it updates QoS Available on behalf of the entire RMD domain. If it cannot (since the M flag is not set for <Path Latency>), it raises the parameter-specific, Not Supported N flag, warning the QNR that the final latency value in QoS Available is imprecise.

In the XG domain, the Initiator QSPEC is translated into a local QSPEC using a similar procedure as described above. The Local QSPEC becomes the current QSPEC used within the XG domain, and the Initiator QSPEC is encapsulated. This saves the QNEs within the XG domain the trouble of re-translating the Initiator QSPEC, and simplifies processing in the local domain. At the egress edge of the XG domain, the translated Local QSPEC is removed, and the Initiator QSPEC returns to the number one position.

If the reservation was successful, eventually the RESERVE request arrives at the QNR (otherwise, the QNE at which the reservation failed aborts the RESERVE and sends an error RESPONSE back to the QNI). If the RII was included in the QoS NSLP message, the QNR generates a positive RESPONSE with QSPEC objects QoS Reserved and QoS

Ash, et al. Experimental

[Page 19]

Available. The parameters appearing in QoS Reserved are the same as in QoS Desired, with values copied from QoS Available. Hence, the QNR includes the following QSPEC objects in the RESPONSE:

QoS Reserved = <TMOD> <PHB Class> QoS Available = <TMOD> <Path Latency>

If the handheld device on the right of Figure 2 is mobile, and moves through different XG wireless networks, then the QoS might change on the path since different XG wireless networks might support different QOSMs. As a result, QoS NSLP/QSPEC processing will have to renegotiate the QoS Available on the path. From a QSPEC perspective, this is like a new reservation on the new section of the path and is basically the same as any other rerouting event -- to the QNEs on the new path, it looks like a new reservation. That is, in this mobile scenario, the new segment may support a different QOSM than the old segment, and the QNI would now signal a new reservation explicitly (or implicitly with the next refreshing RESERVE message) to account for the different QOSM in the XG wireless domain. Further details on rerouting are specified in [RFC5974].

For bit-level examples of QSPECs, see the documents specifying QOSMs: [CL-QOSM], [RFC5976], and [RFC5977].

4. QSPEC Processing and Procedures

Three flags are used in QSPEC processing, the M flag, E flag, and N flag, which are explained in this section. The QNI sets the M flag for each QSPEC parameter it populates that MUST be interpreted by downstream QNEs. If a QNE does not support the parameter, it sets the N flag and fails the reservation. If the QNE supports the parameter but cannot meet the resources requested by the parameter, it sets the E flag and fails the reservation.

If the M flag is not set, the downstream QNE SHOULD interpret the parameter. If the QNE does not support the parameter, it sets the N flag and forwards the reservation. If the QNE supports the parameter but cannot meet the resources requested by the parameter, it sets the E flag and fails the reservation.

4.1. Local QSPEC Definition and Processing

A QNE at the edge of a local domain may either a) translate the Initiator QSPEC into a Local QSPEC and encapsulate the Initiator QSPEC in the RESERVE message, or b) 'hide' the Initiator QSPEC through the local domain and reserve resources by generating a new

Ash, et al. Experimental

[Page 20]

RESERVE message through the local domain containing the Local QSPEC. In either case, the Initiator QSPEC parameters are interpreted at the local domain edges.

A Local QSPEC may allow a simpler control plane in a local domain. The edge nodes in the local domain must interpret the Initiator QSPEC parameters. They can either initiate a parallel session with Local QSPEC or define a Local QSPEC and encapsulate the Initiator QSPEC, as illustrated in Figure 3. The Initiator/Local QSPEC bit identifies whether the QSPEC is an Initiator QSPEC or a Local QSPEC. The QSPEC Type indicates, for example, that the initiator of the local QSPEC uses to a certain QOSM, e.g., CL-QSPEC Type. It may be useful for the QNI to signal a QSPEC Type based on some QOSM (which will necessarily entail populating certain QOSM-related parameters) so that a downstream QNE can chose amongst various QOSM-related processes it might have. That is, the QNI populates the QSPEC Type, e.g., CL-QSPEC Type and sets the Initiator/Local QSPEC bit to 'Initiator'. A local QNE can decide, for whatever reasons, to insert a Local QSPEC Type, e.g., RMD-QSPEC Type, and set the local QSPEC Type = RMD-QSPEC and set the Initiator/Local QSPEC bit to 'Local' (and encapsulate the Initiator QSPEC in the RESERVE or whatever NSLP message).



Figure 3: Defining a Local QSPEC

Ash, et al. Experimental

[Page 21]

Here the QoS-NSLP only sees and passes one QSPEC up to the RMF. Thus, the type of the QSPEC may change within a local domain. Hence:

- o the QNI signals its QoS requirements with the Initiator QSPEC,
- o the ingress edge QNE in the local domain translates the Initiator QSPEC parameters to equivalent parameters in the local QSPEC,
- o the QNEs in the local domain only interpret the Local QSPEC parameters, and
- o the egress QNE in the local domain processes the Local QSPEC and also interprets the QSPEC parameters in the Initiator QSPEC.

The Local QSPEC MUST be consistent with the Initiator QSPEC. That is, it MUST NOT specify a lower level of resources than specified by the Initiator QSPEC. For example, in RMD the TMOD parameters contained in the original Initiator QSPEC are mapped to the equivalent TMOD parameter representing only the peak bandwidth in the Local QSPEC.

Note that it is possible to use both a) hiding a QSPEC through a local domain by initiating a new RESERVE at the domain edge, and b) defining a Local QSPEC and encapsulating the Initiator QSPEC, as defined above. However, it is not expected that both the hiding and encapsulating functions would be used at the same time for the same flow.

The support of Local QSPECs is illustrated in Figure 4 for a single flow to show where the Initiator and Local QSPECs are used. The QNI initiates an end-to-end, inter-domain QoS NSLP RESERVE message containing the Initiator QSPEC for the Y.1541 QOSM. As illustrated in Figure 4, the RESERVE message crosses multiple domains supporting different QOSMs. In this illustration, the Initiator QSPEC arrives in a QoS NSLP RESERVE message at the ingress node of the local-QOSM domain. At the ingress edge node of the local-QOSM domain, the endto-end, inter-domain QoS-NSLP message triggers the generation of a Local QSPEC, and the Initiator QSPEC is encapsulated within the messages signaled through the local domain. The local QSPEC is used for QoS processing in the local-QOSM domain, and the Initiator QSPEC is used for QoS processing outside the local domain.

In this example, the QNI sets <QoS Desired>, <Minimum QoS>, and <QoS Available> objects to include objectives for the <Path Latency>, <Path Jitter>, and <Path PER> parameters. The QNE / local domain sets the cumulative parameters, e.g., <Path Latency>, that can be achieved in the <QoS Available> object (but not less than specified in <Minimum QoS>). If the <QoS Available> fails to satisfy one or

Ash, et al. Experimental

[Page 22]

more of the <Minimum QoS> objectives, the QNE / local domain notifies the QNI and the reservation is aborted. If any QNE cannot meet the requirements designated by the Initiator QSPEC to support a QSPEC parameter with the M bit set to zero, the QNE sets the N flag for that parameter to one. Otherwise, the QNR notifies the QNI of the <QoS Available> for the reservation.

j	e2e	<->	e2e	<				>	e2e	<->	e2e	
ĺ	QOSM		QOSM						QOSM		QOSM	
ĺ												
Í	NSLP		NSLP	<->	NSLP	<->	NSLP	<->	NSLP		NSLP	
ĺ	Y.1541		local		local		local		local		Y.1541	
	QOSM		QOSM		QOSM		QOSM		QOSM		QOSM	
-												-
	NTLP	<->	NTLP	<->	NTLP	<->	NTLP	<->	NTLP	<->	NTLP	
	QNI		QNE		QNE		QNE		QNE		QNR	
((End)	(Ingi	cess Edg	ge) (Interio	<u>(</u>)	Interio	_) (I	Igress H	Edge () (End)	1

Figure 4: Example of Initiator and Local Domain QOSM Operation

4.2. Reservation Success/Failure, QSPEC Error Codes, and INFO-SPEC Notification

A reservation may not be successful for several reasons:

- a reservation may fail because the desired resources are not available. This is a reservation failure condition.
- a reservation may fail because the QSPEC is erroneous or because of a QNE fault. This is an error condition.

A reservation may be successful even though some parameters could not be interpreted or updated properly:

- a QSPEC parameter cannot be interpreted because it is an unknown QSPEC parameter type. This is a QSPEC parameter not supported condition. However, the reservation does not fail. The QNI can still decide whether to keep or tear down the reservation depending on the procedures specified by the QNI's QOSM.

The following sections provide details on the handling of unsuccessful reservations and reservations where some parameters could not be met, as follows:

Ash, et al. Experimental

[Page 23]

- details on flags used inside the QSPEC to convey information on success or failure of individual parameters. The formats and semantics of all flags are given in Section 5.
- the content of the INFO-SPEC [RFC5974], which carries a code indicating the outcome of reservations.
- the generation of a RESPONSE message to the QNI containing both QSPEC and INFO-SPEC objects.

Note that when there are routers along the path between the QNI and QNR where QoS cannot be provided, then the QoS-NSLP generic flag BREAK (B) is set. The BREAK flag is discussed in Section 3.3.5 of [RFC5974].

4.2.1. Reservation Failure and Error E Flag

The QSPEC parameters each have a 'reservation failure error E flag' to indicate which (if any) parameters could not be satisfied. When a resource cannot be satisfied for a particular parameter, the QNE detecting the problem raises the E flag in this parameter. Note that the TMOD parameter and all QSPEC parameters with the M flag set MUST be examined by the RMF, and all QSPEC parameters with the M flag not set SHOULD be examined by the RMF, and the E flag set to indicate whether the parameter could or could not be satisfied. Additionally, the E flag in the corresponding QSPEC object MUST be raised when a resource cannot be satisfied for this parameter. If the reservation failure problem cannot be located at the parameter level, only the E flag in the QSPEC object is raised.

When an RMF cannot interpret the QSPEC because the coding is erroneous, it raises corresponding reservation failure E flags in the QSPEC. Normally, all QSPEC parameters MUST be examined by the RMF, and the erroneous parameters appropriately flagged. In some cases, however, an error condition may occur and the E flag of the errorcausing QSPEC parameter is raised (if possible), but the processing of further parameters may be aborted.

Note that if the QSPEC and/or any QSPEC parameter is found to be erroneous, then any QSPEC parameters not satisfied are ignored and the E Flags in the QSPEC object MUST NOT be set for those parameters (unless they are erroneous).

Whether E flags denote reservation failure or error can be determined by the corresponding error code in the INFO-SPEC in QoS NSLP, as discussed below.

Ash, et al. Experimental

[Page 24]

RFC 5975

4.2.2. QSPEC Parameter Not Supported N Flag

Each QSPEC parameter has an associated 'Not Supported N flag'. If the Not Supported N flag is set, then at least one QNE along the data transmission path between the QNI and QNR cannot interpret the specified QSPEC parameter. A QNE MUST set the Not Supported N flag if it cannot interpret the QSPEC parameter. If the M flag for the parameter is not set, the message should continue to be forwarded but with the N flag set, and the QNI has the option of tearing down the reservation.

If a QNE in the path does not support a QSPEC parameter, e.g., <Path Latency>, and sets the N flag, then downstream QNEs that support the parameter SHOULD still update the parameter, even if the N flag is set. However, the presence of the N flag will indicate that the cumulative value only provides a bound, and the QNI/QNR decides whether or not to accept the reservation with the N flag set.

4.2.3. INFO-SPEC Coding of Reservation Outcome

As prescribed by [RFC5974], the RESPONSE message always contains the INFO-SPEC with an appropriate 'error' code. It usually also contains a QSPEC with QSPEC objects, as described in Section 4.3 ("QSPEC Procedures"). The RESPONSE message MAY omit the QSPEC in case of a successful reservation.

The following guidelines are provided for setting the error codes in the INFO-SPEC, based on the codes provided in Section 5.1.3.6 of [RFC5974]:

- NSLP error class 2 (Success) / 0x01 (Reservation Success): This code is set when all QSPEC parameters have been satisfied. In this case, no E Flag is set; however, one or more N flags may be set.
- NSLP error class 4 (Transient Failure) / 0x07 (Reservation Failure): This code is set when at least one QSPEC parameter could not be satisfied, or when a QSPEC parameter with M flag set could not be interpreted. E flags are set for the parameters that could not be satisfied at each QNE up to the QNE issuing the RESPONSE message. The N flag is set for those parameters that could not be interpreted by at least one QNE. In this case, QNEs receiving the RESPONSE message MUST remove the corresponding reservation.

Ash, et al. Experimental

[Page 25]

- NSLP error class 3 (Protocol Error) / 0x0c (Malformed QSPEC): Some QSPEC parameters had associated errors, E Flags are set for parameters that had errors, and the QNE where the error was found rejects the reservation.
- NSLP error class 3 (Protocol Error) / 0x0f (Incompatible QSPEC): A higher version QSPEC is signaled and not supported by the QNE.
- NSLP error class 6 (QoS Model Error): QOSM error codes can be defined by QOSM specification documents. A registry is defined in Section 7, IANA Considerations.
- 4.2.4. QNE Generation of a RESPONSE Message
 - Successful Reservation Condition

When a RESERVE message arrives at a QNR and no E Flag is set, the reservation is successful. A RESPONSE message may be generated with INFO-SPEC code 'Reservation Success' as described above and in Section 4.3 ("QSPEC Procedures").

- Reservation Failure Condition

When a QNE detects that a reservation failure occurs for at least one parameter, the QNE sets the E Flags for the QSPEC parameters and QSPEC object that failed to be satisfied. According to [RFC5974], the QNE behavior depends on whether it is stateful or not. When a stateful QNE determines the reservation failed, it formulates a RESPONSE message that includes an INFO-SPEC with the 'reservation failure' error code and QSPEC object. The QSPEC in the RESPONSE message includes the failed QSPEC parameters marked with the E Flag to clearly identify them.

The default action for a stateless QoS NSLP QNE that detects a reservation failure condition is that it MUST continue to forward the RESERVE message to the next stateful QNE, with the E Flags appropriately set for each QSPEC parameter. The next stateful QNE then formulates the RESPONSE message as described above.

- Malformed QSPEC Error Condition

When a stateful QNE detects that one or more QSPEC parameters are erroneous, the QNE sets the error code 'malformed QSPEC' in the INFO-SPEC. In this case, the QSPEC object with the E Flags appropriately set for the erroneous parameters is returned within the INFO-SPEC object. The QSPEC object can be truncated or fully included within the INFO-SPEC.

Ash, et al. Experimental

[Page 26]

According to [RFC5974], the QNE behavior depends on whether it is stateful or not. When a stateful QNE determines a malformed QSPEC error condition, it formulates a RESPONSE message that includes an INFO-SPEC with the 'malformed QSPEC' error code and QSPEC object.

The QSPEC in the RESPONSE message includes, if possible, only the erroneous QSPEC parameters and no others. The erroneous QSPEC parameter(s) are marked with the E Flag to clearly identify them. If QSPEC parameters are returned in the INFO-SPEC that are not marked with the E flag, then any values of these parameters are irrelevant and MUST be ignored by the QNI.

The default action for a stateless QoS NSLP QNE that detects a malformed QSPEC error condition is that it MUST continue to forward the RESERVE message to the next stateful QNE, with the E Flags appropriately set for each QSPEC parameter. The next stateful QNE will then act as described in [RFC5974].

A 'malformed QSPEC' error code takes precedence over the 'reservation failure' error code, and therefore the case of reservation failure and QSPEC/RMF error conditions are disjoint, and the same E Flag can be used in both cases without ambiguity.

4.2.5. Special Case of Local QSPEC

When an unsuccessful reservation problem occurs inside a local domain where a Local QSPEC is used, only the topmost (local) QSPEC is affected (e.g., E flags are raised, etc.). The encapsulated Initiator QSPEC is untouched. However, when the message (RESPONSE in case of stateful QNEs; RESERVE in case of stateless QNEs) reaches the edge of the local domain, the Local QSPEC is removed. The edge QNE must update the Initiator QSPEC on behalf of the entire domain, reflecting the information received in the Local OSPEC. This update concerns both parameter values and flags. Note that some intelligence is needed in mapping the E flags, etc., from the local QSPEC to the Initiator QSPEC. For example, even if there is no direct match between the parameters in the local and Initiator QSPECs, E flags could still be raised in the latter.

4.3. QSPEC Procedures

While the QSPEC template aims to put minimal restrictions on usage of QSPEC objects, interoperability between QNEs and between QOSMs must be ensured. We therefore give below an exhaustive list of QSPEC object combinations for the message sequences described in QoS NSLP [RFC5974]. A specific QOSM may prescribe that only a subset of the procedures listed below may be used.

Ash, et al. Experimental

[Page 27]

Note that QoS NSLP does not mandate the usage of a RESPONSE message. A positive RESPONSE message will only be generated if the QNE includes an RII (Request Identification Information) in the RESERVE message, and a negative RESPONSE message is always generated in case of an error or failure. Some of the QSPEC procedures below, however, are only meaningful when a RESPONSE message is possible. The QNI SHOULD in these cases include an RII.

4.3.1. Two-Way Transactions

Here, the QNI issues a RESERVE message, which may be replied to by a RESPONSE message. The following 3 cases for QSPEC object usage exist:

MESSAGE SEQUENCE	OBJECT COMBINATION	OBJECTS INCLUDED IN RESERVE MESSAGE	OBJECTS INCLUDED IN RESPONSE MESSAGE
0	0	QoS Desired	QoS Reserved
0	1	QoS Desired QoS Available	QoS Reserved QoS Available
0	2	QoS Desired QoS Available Minimum QoS	QoS Reserved QoS Available

Table 1: Message Sequence 0: Two-Way Transactions Defining Object Combinations 0, 1, and 2

Case 1:

If only QoS Desired is included in the RESERVE message, the implicit assumption is that exactly these resources must be reserved. If this is not possible, the reservation fails. The parameters in QoS Reserved are copied from the parameters in QoS Desired. If the reservation is successful, the RESPONSE message can be omitted in this case. If a RESPONSE message was requested by a QNE on the path, the QSPEC in the RESPONSE message can be omitted.

Case 2:

When QoS Available is included in the RESERVE message also, some parameters will appear only in QoS Available and not in QoS Desired. It is assumed that the value of these parameters is collected for informational purposes only (e.g., path latency).

Ash, et al. Experimental

[Page 28]

However, some parameters in QoS Available can be the same as in QoS Desired. For these parameters, the implicit message is that the QNI would be satisfied by a reservation with lower parameter values than specified in QoS Desired. For these parameters, the QNI seeds the parameter values in QoS Available to those in QoS Desired (except for cumulative parameters such as <Path Latency>).

Each QNE interprets the parameters in QoS Available according to its current capabilities. Reservations in each QNE are hence based on current parameter values in QoS Available (and additionally those parameters that only appear in QoS Desired). The drawback of this approach is that, if the resulting resource reservation becomes gradually smaller towards the QNR, QNEs close to the QNI have an oversized reservation, possibly resulting in unnecessary costs for the user. Of course, in the RESPONSE the QNI learns what the actual reservation is (from the QoS RESERVED object) and can immediately issue a properly sized refreshing RESERVE. The advantage of the approach is that the reservation is performed in half-a-roundtrip time.

The QSPEC parameter IDs and values included in the QoS Reserved object in the RESPONSE message MUST be the same as those in the QoS Desired object in the RESERVE message. For those QSPEC parameters that were also included in the QoS Available object in the RESERVE message, their value is copied from the QoS Available object (in RESERVE) into the QoS Reserved object (in RESPONSE). For the other QSPEC parameters, the value is copied from the QoS Desired object (the reservation would fail if the corresponding QoS could not be reserved).

All parameters in the QoS Available object in the RESPONSE message are copied with their values from the QoS Available object in the RESERVE message (irrespective of whether they have also been copied into the QoS Desired object). Note that the parameters in the QoS Available object can be overwritten in the RESERVE message, whereas they cannot be overwritten in the RESPONSE message.

In this case, the QNI SHOULD request a RESPONSE message since it will otherwise not learn what QoS is available.

Case 3:

This case is handled as case 2, except that the reservation fails when QoS Available becomes less than Minimum QoS for one parameter. If a parameter appears in the QoS Available object but not in the Minimum QoS object, it is assumed that there is no minimum value for this parameter.

Ash, et al. Experimental

[Page 29]

Regarding Traffic Handling Directives, the default rule is that all QSPEC parameters that have been included in the RESERVE message by the QNI are also included in the RESPONSE message by the QNR with the value they had when arriving at the QNR. When traveling in the RESPONSE message, all Traffic Handling Directives parameters are read-only. Note that a QOSM specification may define its own Traffic Handling Directives parameters and processing rules.

4.3.2. Three-Way Transactions

Here, the QNR issues a QUERY message that is replied to by the QNI with a RESERVE message if the reservation was successful. The QNR in turn sends a RESPONSE message to the QNI. The following 3 cases for QSPEC object usage exist:

MSG. |OBJ. |OBJECTS INCLUDED |OBJECTS INCLUDED |OBJECTS INCLUDED SEQ. COM. IN QUERY MESSAGE IN RESERVE MESSAGE IN RESPONSE MESSAGE _____

1	0	QoS Desired	QoS Desired	QoS Reserved
1	1	QoS Desired (Minimum QoS)	QoS Desired QoS Available (Minimum QoS)	QoS Reserved QoS Available
1	2	 QoS Desired QoS Available	 QoS Desired QoS Available	QoS Reserved

Table 2: Message Sequence 1: Three-Way Transactions Defining Object Combinations 0, 1, and 2

Cases 1 and 2:

The idea is that the sender (QNR in this scenario) needs to inform the receiver (QNI in this scenario) about the QoS it desires. To this end, the sender sends a QUERY message to the receiver including a QoS Desired QSPEC object. If the QoS is negotiable, it additionally includes a (possibly zero) Minimum QoS object, as in Case 2.

The RESERVE message includes the QoS Available object if the sender signaled that QoS is negotiable (i.e., it included the Minimum QoS object). If the Minimum QoS object received from the sender is included in the QUERY message, the QNI also includes the Minimum QoS object in the RESERVE message.

Ash, et al. Experimental

[Page 30]

For a successful reservation, the RESPONSE message in case 1 is optional (as is the QSPEC inside). In case 2, however, the RESPONSE message is necessary in order for the QNI to learn about the QoS available.

Case 3:

This is the 'RSVP-style' scenario. The sender (QNR in this scenario) issues a QUERY message with a QoS Desired object informing the receiver (QNI in this scenario) about the QoS it desires, as above. It also includes a QoS Available object to collect path properties. Note that here path properties are collected with the QUERY message, whereas in the previous case, 2 path properties were collected in the RESERVE message.

Some parameters in the QoS Available object may be the same as in the QoS Desired object. For these parameters, the implicit message is that the sender would be satisfied by a reservation with lower parameter values than specified in QoS Desired.

It is possible for the QoS Available object to contain parameters that do not appear in the QoS Desired object. It is assumed that the value of these parameters is collected for informational purposes only (e.g., path latency). Parameter values in the QoS Available object are seeded according to the sender's capabilities. Each QNE remaps or approximately interprets the parameter values according to its current capabilities.

The receiver (QNI in this scenario) signals the QoS Desired object as follows: For those parameters that appear in both the QoS Available object and QoS Desired object in the QUERY message, it takes the (possibly remapped) QSPEC parameter values from the QoS Available object. For those parameters that only appear in the QoS Desired object, it adopts the parameter values from the QoS Desired object.

The parameters in the QoS Available QSPEC object in the RESERVE message are copied with their values from the QoS Available QSPEC object in the QUERY message. Note that the parameters in the QoS Available object can be overwritten in the QUERY message, whereas they cannot be overwritten in the RESERVE message.

The advantage of this model compared to the sender-initiated reservation is that the situation of over-reservation in QNEs close to the QNI (as described above) does not occur. On the other hand, the QUERY message may find, for example, a particular bandwidth is

Ash, et al. Experimental

[Page 31]

not available. When the actual reservation is performed, however, the desired bandwidth may meanwhile have become free. That is, the 'RSVP style' may result in a smaller reservation than necessary.

The sender includes all QSPEC parameters it cares about in the QUERY message. Parameters that can be overwritten are updated by QNEs as the QUERY message travels towards the receiver. The receiver includes all QSPEC parameters arriving in the QUERY message also in the RESERVE message, with the value they had when arriving at the receiver. Again, QOSM-specific QSPEC parameters and procedures may be defined in QOSM specification documents.

Also in this scenario, the QNI SHOULD request a RESPONSE message since it will otherwise not learn what QoS is available.

Regarding Traffic Handling Directives, the default rule is that all QSPEC parameters that have been included in the RESERVE message by the QNI are also included in the RESPONSE message by the QNR with the value they had when arriving at the QNR. When traveling in the RESPONSE message, all Traffic Handling Directives parameters are read-only. Note that a QOSM specification may define its own Traffic Handling Directives parameters and processing rules.

4.3.3. Resource Queries

Here, the QNI issues a QUERY message in order to investigate what resources are currently available. The QNR replies with a RESPONSE message.

MESSAGE		OBJECTS INCLUDED	OBJECTS INCLUDED
SEQUENCE		IN QUERY MESSAGE	IN RESPONSE MESSAGE
2	0	QoS Available	QoS Available

Table 3: Message Sequence 2: Resource Queries Defining Object Combination 0

Note that the QoS Available object when traveling in the QUERY message can be overwritten, whereas in the RESPONSE message it cannot be overwritten.

Regarding Traffic Handling Directives, the default rule is that all QSPEC parameters that have been included in the RESERVE message by the QNI are also included in the RESPONSE message by the QNR with the value they had when arriving at the QNR. When traveling in the RESPONSE message, all Traffic Handling Directives parameters are read-only. Note that a QOSM specification may define its own Traffic Handling Directives parameters and processing rules.

Ash, et al. Experimental

[Page 32]

4.3.4. Bidirectional Reservations

On a QSPEC level, bidirectional reservations are no different from unidirectional reservations, since QSPECs for different directions never travel in the same message.

4.3.5. Preemption

A flow can be preempted by a QNE based on QNE policy, where a decision to preempt a flow may account for various factors such as, for example, the values of the QSPEC preemption priority and defending priority parameters as described in Section 5.2.8. In this case, the reservation state for this flow is torn down in the QNE, and the QNE sends a NOTIFY message to the QNI, as described in [RFC5974]. The NOTIFY message carries an INFO-SPEC with the error code as described in [RFC5974]. A QOSM specification document may specify whether a NOTIFY message also carries a QSPEC object. The QNI would normally tear down the preempted reservation by sending a RESERVE message with the TEAR flag set using the SII of the preempted reservation. However, the QNI can follow other procedures as specified in its QOSM specification document.

4.4. QSPEC Extensibility

Additional QSPEC parameters MAY need to be defined in the future and are defined in separate informational documents. For example, QSPEC parameters are defined in [RFC5977] and [RFC5976].

Guidelines on the technical criteria to be followed in evaluating requests for new codepoint assignments for QSPEC objects and QSPEC parameters are given in Section 7, IANA Considerations.

5. QSPEC Functional Specification

This section defines the encodings of the QSPEC parameters. We first give the general QSPEC formats and then the formats of the QSPEC objects and parameters.

Network octet order ('big-endian') for all 16- and 32-bit integers, as well as 32-bit floating point numbers, is as specified in [RFC4506], [IEEE754], and [NETWORK-OCTET-ORDER].

5.1. General QSPEC Formats

The format of the QSPEC closely follows that used in GIST [RFC5971] and QoS NSLP [RFC5974]. Every object (and parameter) has the following general format:

Ash, et al. Experimental

[Page 33]

- o The overall format is Type-Length-Value (in that order).
- o Some parts of the type field are set aside for control flags.
- o Length has the units of 32-bit words, and measures the length of Value. If there is no Value, Length=0. The Object length excludes the header.
- o Value is a whole number of 32-bit words. If there is any padding required, the length and location MUST be defined by the objectspecific format information; objects that contain variable-length types may need to include additional length subfields to do so.
- o Any part of the object used for padding or defined as reserved ("r") MUST be set to 0 on transmission and MUST be ignored on reception.
- o Empty QSPECs and empty QSPEC Objects MUST NOT be used.
- o Duplicate objects, duplicate parameters, and/or multiple occurrences of a parameter MUST NOT be used.

0	1	2	3									
0 1 2 3 4 5 6	7 8 9 0 1 2 3 4	5 6 7 8 9 0 1 2 3 4	5678901									
+-												
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 +-+++++++++++++++++++++++++++++++++++												
~ 1												
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-												
+-+-+-+-+-+-	+-	-+	+-+-+-+-+-+-+									

5.1.1. Common Header Format

The Common QSPEC Header is a fixed 4-octet object containing the QSPEC Version, QSPEC Type, an identifier for the QSPEC Procedure (see Section 4.3), and an Initiator/Local QSPEC bit:

0									1										2										3	
0	1 2	3	4	5	б	7	8	9	0	1	2	3	4	5	б	7	8	9	0	1	2	3	4	5	б	7	8	9	0	1
+-+	-+-	+ - +	+	+	⊢	+ - +	⊢ – +	+	+ - +	⊢	+	+	+ - +	⊢	+	+	+	+	+	+	+	+	+	+	+ - •	+	+	+	+	⊦-+
+-+-+-+-+-+-+-+-+-+-+-++-++-++-++-++-++																														
+-+	-+-	+ - +	+	+	⊦ — -	+ - +	+ +	+	+ - +	⊦ — -	+	+	+ - +	⊦ — -	+	+	+	+	+	+	+	+	+	+	+ - •	+	+	+	+	+-+

Vers.: Identifies the QSPEC version number. QSPEC Version 0 is assigned by this specification in Section 7 (IANA Considerations).

Ash, et al. Experimental

[Page 34]

- QSPEC Type: Identifies the particular type of QSPEC, e.g., a QSPEC Type corresponding to a particular QOSM. QSPEC Type 0 (default) is assigned by this specification in Section 7 (IANA Considerations).
- QSPEC Proc.: Identifies the QSPEC procedure and is composed of two times 4 bits. The first field identifies the Message Sequence; the second field identifies the QSPEC Object Combination used for this particular message sequence:

0 1 2 3 4 5 6 7 |Mes.Sq |Obj.Cmb| +-+-+-+-+-+-+-+-+

The Message Sequence field can attain the following values:

0: Sender-Initiated Reservations 1: Receiver-Initiated Reservations 2: Resource Queries

The Object Combination field can take the values between 1 and 3 indicated in the tables in Section 4.3:

Message Sequence: 0 Object Combination: 0, 1, 2 Semantic: see Table 1 in Section 4.3.1

Message Sequence: 1 Object Combination: 0, 1, 2 Semantic: see Table 2 in Section 4.3.2

Message Sequence: 2 Object Combination: 0 Semantic: see Table 3 in Section 4.3.3

I: Initiator/Local QSPEC bit identifies whether the QSPEC is an initiator QSPEC or a Local QSPEC, and is set to the following values:

> 0: Initiator QSPEC 1: Local QSPEC

Length: The total length of the QSPEC (in 32-bit words) excluding the common header

Ash, et al.

Experimental

[Page 35]

The QSPEC Objects field is a collection of QSPEC objects (QoS Desired, QoS Available, etc.), which share a common format and each contain several parameters.

5.1.2. QSPEC Object Header Format

QSPEC objects share a common header format:

0 1 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 |E|r|r|r| Object Type |r|r|r| Length | E Flag: Set if an error occurs on object level Object Type = 0: QoS Desired (parameters cannot be overwritten) = 1: QoS Available (parameters may be overwritten; see Section 3.2) = 2: QoS Reserved (parameters cannot be overwritten) = 3: Minimum QoS (parameters cannot be overwritten) The r bits are reserved. Each QSPEC or QSPEC parameter within an object is encoded in the same way in TLV format using a similar parameter header: 1 2 3 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 |M|E|N|r| Parameter ID |r|r|r|r| Length M Flag: When set, indicates the subsequent parameter MUST be interpreted. Otherwise, the parameter can be ignored if not understood. E Flag: When set, indicates either a) a reservation failure where the QSPEC parameter is not met, or b) an error occurred when this parameter was being interpreted (see Section 4.2.1).

N Flag: Not Supported QSPEC parameter flag (see Section 4.2.2).

Parameter ID: Assigned consecutively to each QSPEC parameter. Parameter IDs are assigned to each QSPEC parameter defined in this document in Sections 5.2 and 7 (IANA Considerations).

Ash, et al. Experimental [Page 36]
Parameters are usually coded individually, for example, the <Excess Treatment> parameter (Section 5.2.11). However, it is also possible to combine several sub-parameters into one parameter field, which is called 'container coding'. This coding is useful if either a) the sub-parameters always occur together (as for example the 5 subparameters that jointly make up the TMOD), or b) in order to make coding more efficient when the length of each sub-parameter value is much less than a 32-bit word (as for example described in [RFC5977]) and to avoid header overload. When a container is defined, the Parameter ID and the M, E, and N flags refer to the container. Examples of container parameters are <TMOD> (specified below) and the PHR (Per Hop Reservation) container parameter specified in [RFC5977].

5.2. QSPEC Parameter Coding

The references in the following sections point to the normative procedures for processing the QSPEC parameters and sub-parameters.

5.2.1. <TMOD-1> Parameter

The <TMOD-1> parameter consists of the <r>, , , <m>, and <MPS> sub-parameters [RFC2212], which all must be populated in the <TMOD-1> parameter. Note that a second TMOD QSPEC parameter <TMOD-2> is specified below in Section 5.2.2.

The coding for the <TMOD-1> parameter is as follows:

1 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 |1|E|0|r| 1 |r|r|r|r| 5 TMOD Rate-1 (r) (32-bit IEEE floating point number) TMOD Size-1 (b) (32-bit IEEE floating point number) Peak Data Rate-1 (p) (32-bit IEEE floating point number) Minimum Policed Unit-1 (m) (32-bit unsigned integer) Maximum Packet Size-1 (MPS) (32-bit unsigned integer)

The <TMOD-1> parameters are represented by three floating point numbers in single-precision IEEE floating point format [IEEE754] followed by two 32-bit integers in network octet order. The first floating point value is the rate (r), the second floating point value is the bucket size (b), the third floating point is the peak rate

Ash, et al. Experimental

[Page 37]

(p), the first unsigned integer is the minimum policed unit (m), and the second unsigned integer is the maximum packet size (MPS). The values of r and p are measured in octets per second; b, m, and MPS are measured in octets. When r, b, and p terms are represented as IEEE floating point values, the sign bit MUST be zero (all values MUST be non-negative). Exponents less than 127 (i.e., 0) are prohibited. Exponents greater than 162 (i.e., positive 35) are discouraged, except for specifying a peak rate of infinity. Infinity is represented with an exponent of all ones (255), and a sign bit and mantissa of all zeroes.

5.2.2. <TMOD-2> Parameter

A second QSPEC <TMOD-2> parameter is specified as could be needed, for example, to support some Diffserv applications.

The coding for the <TMOD-2> parameter is as follows:

0 1 2 3 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 |M|E|N|r| 2 |r|r|r|r| 5| TMOD Rate-2 (r) (32-bit IEEE floating point number) | TMOD Size-2 (b) (32-bit IEEE floating point number) Peak Data Rate-2 (p) (32-bit IEEE floating point number) | Minimum Policed Unit-2 (m) (32-bit unsigned integer) Maximum Packet Size-2 (MPS) (32-bit unsigned integer)

The <TMOD-2> parameters are represented by three floating point numbers in single-precision IEEE floating point format [IEEE754] followed by two 32-bit integers in network octet order. The first floating point value is the rate (r), the second floating point value is the bucket size (b), the third floating point is the peak rate (p), the first unsigned integer is the minimum policed unit (m), and the second unsigned integer is the maximum packet size (MPS). The values of r and p are measured in octets per second; b, m, and MPS are measured in octets. When r, b, and p terms are represented as IEEE floating point values, the sign bit MUST be zero (all values MUST be non-negative). Exponents less than 127 (i.e., 0) are prohibited. Exponents greater than 162 (i.e., positive 35) are

Ash, et al. Experimental

[Page 38]

discouraged, except for specifying a peak rate of infinity. Infinity is represented with an exponent of all ones (255), and a sign bit and mantissa of all zeroes.

5.2.3. <Path Latency> Parameter

0	1	2	3
0 1 2 3 4 5 6	7 8 9 0 1 2 3	4 5 6 7 8 9 0 1 2 3	45678901
+-+-+-+-+-+-+	-+-+-+-+-+-+	+-	-+-+-+-+-+-+-+-+
M E N r	3	r r r r	1
++++	-++	+++++	++++
	Path Latency	y (32-bit unsigned in	teger)
+-+-+-+-+-+-+	-+-+-+-+-+-+	+-	-+-+-+-+-+-+-+-+

The Path Latency [RFC2215] is a single 32-bit unsigned integer in network octet order. The intention of the Path Latency parameter is the same as the Minimal Path Latency parameter defined in Section 3.4 of [RFC2215]. The purpose of this parameter is to provide a baseline minimum path latency for use with services that provide estimates or bounds on additional path delay, such as in [RFC2212]. Together with the queuing delay bound offered by [RFC2212] and similar services, this parameter gives the application knowledge of both the minimum and maximum packet delivery delay.

The composition rule for the <Path Latency> parameter is summation with a clamp of $(2^{32}) - 1$ on the maximum value. The latencies are average values reported in units of one microsecond. A system with resolution less than one microsecond MUST set unused digits to zero. An individual QNE can add a latency value between 1 and 2^28 (somewhat over two minutes), and the total latency added across all QNEs can range as high as $(2^{32})-2$. If the sum of the different elements delays exceeds $(2^32)-2$, the end-to-end cumulative delay SHOULD be reported as indeterminate = $(2^{32})-1$. A QNE that cannot accurately predict the latency of packets it is processing MUST raise the Not Supported N flag and either leave the value of Path Latency as is, or add its best estimate of its lower bound. A raised notsupported flag indicates the value of Path Latency is a lower bound of the real Path Latency. The distinguished value $(2^32)-1$ is taken to mean indeterminate latency because the composition function limits the composed sum to this value; it indicates the range of the composition calculation was exceeded.

Ash, et al. Experimental

[Page 39]

5.2.4. <Path Jitter> Parameter

			23456789	
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-	4	r r r r	+-	+-+-+
+-+-+-+-+-++++ Path Jitte +-+-+-+-+-++++ Path Jitter S +-+-+++++++++++ Path Jitte	r STAT1(varian -+-+-+-+-+-+-+ r STAT2(99.9%- -+-+-+-+-+-+ TAT3(minimum L -+-+-+++++++++ r STAT4(Reserv	ile) (32-bit uns -+-+-+-+-+-+	+-+-+-+-+-+-+-+	 +-+-+ (-) +-+-+ (-)

The Path Jitter is a set of four 32-bit unsigned integers in network octet order [RFC3393, Y.1540, Y.1541]. As noted in Section 3.3.2, the Path Jitter parameter is called "IP Delay Variation" in [RFC3393]. The Path Jitter parameter is the combination of four statistics describing the Jitter distribution with a clamp of (2^32) - 1 on the maximum of each value. The jitter STATs are reported in units of one microsecond. A system with resolution less than one microsecond MUST set unused digits to zero. An individual QNE can add jitter values between 1 and 2^28 (somewhat over two minutes), and the total jitter computed across all QNEs can range as high as $(2^{32})-2$. If the combination of the different element values exceeds $(2^{32})-2$, the end-to-end cumulative jitter SHOULD be reported as indeterminate. A QNE that cannot accurately predict the jitter of packets it is processing MUST raise the not-supported flag and either leave the value of Path Jitter as is, or add its best estimate of its STAT values. A raised not-supported flag indicates the value of Path Jitter is a lower bound of the real Path Jitter. The distinguished value (2^32)-1 is taken to mean indeterminate jitter. A QNE that cannot accurately predict the jitter of packets it is processing SHOULD set its local Path Jitter parameter to this value. Because the composition function limits the total to this value, receipt of this value at a network element or application indicates that the true Path Jitter is not known. This MAY happen because one or more network elements could not supply a value or because the range of the composition calculation was exceeded.

NOTE: The Jitter composition function makes use of the <Path Latency> parameter. Composition functions for loss, latency, and jitter may be found in [Y.1541]. Development continues on methods to combine jitter values to estimate the value of the complete path, and additional statistics may be needed to support new methods (the methods are standardized in [RFC5481] and [COMPOSITION]).

Ash, et al. Experimental

[Page 40]

5.2.5. <Path PLR> Parameter

0	1	2	3
0 1 2 3 4 5	6 7 8 9 0 1 2 3	4 5 6 7 8 9 0 1 2	3 4 5 6 7 8 9 0 1
+-+-+-+-+-+	-+-+-+-+-+-+-+-+	-+	-+
M E N r	5	r r r r	1
++	Path Packet Los	s Ratio (32-bit fl	oating point)

The Path PLR is a single 32-bit single precision IEEE floating point number in network octet order [Y.1541]. As defined in [Y.1540], Path PLR is the ratio of total lost IP packets to total transmitted IP packets. An evaluation interval of 1 minute is suggested in [Y.1541], in which the number of losses observed is directly related to the confidence in the result. The composition rule for the <Path PLR> parameter is summation with a clamp of 10^-1 on the maximum value. The PLRs are reported in units of 10^-11. A system with resolution less than 10^-11 MUST set unused digits to zero. An individual QNE adds its local PLR value (up to a maximum of 10⁻²) to the total Path PLR value (up to a maximum of 10^-1) , where the acceptability of the total Path PLR value added across all QNEs is determined based on the QOSM being used. The maximum limit of 10^-2 on a QNE's local PLR value and the maximum limit (clamp value) of 10^-1 on the accumulated end-to-end Path PLR value are used to preserve the accuracy of the simple additive accumulation function specified and to avoid more complex accumulation functions. Furthermore, if these maximums are exceeded, then the path would likely not meet the QoS objectives. If the sum of the different elements' values exceeds 10^-1, the end-to-end cumulative PLR SHOULD be reported as indeterminate. A QNE that cannot accurately predict the PLR of packets it is processing MUST raise the not-supported flag and either leave the value of Path PLR as is, or add its best estimate of its lower bound. A raised not-supported flag indicates the value of Path PLR is a lower bound of the real Path PLR. The distinguished value 10^-1 is taken to mean indeterminate PLR. A QNE that cannot accurately predict the PLR of packets it is processing SHOULD set its local path PLR parameter to this value. Because the composition function limits the composed sum to this value, receipt of this value at a network element or application indicates that the true path PLR is not known. This MAY happen because one or more network elements could not supply a value or because the range of the composition calculation was exceeded.

Ash, et al. Experimental

[Page 41]

5.2.6. <Path PER> Parameter

0	1	2	3
0 1 2 3 4 5	6 7 8 9 0 1 2 3	4 5 6 7 8 9 0 1 2	345678901
+-+-+-+-+-+	+-+-+-+-+-+-+-	+-	+-
M E N r	6	r r r r	1
	Path Packet Er	ror Ratio (32-bit +-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-	1

The Path PER is a single 32-bit single precision IEEE floating point number in network octet order [Y.1541]. As defined in [Y.1540], Path PER is the ratio of total errored IP packets to the total of successful IP Packets plus errored IP packets, in which the number of errored packets observed is directly related to the confidence in the result. The composition rule for the <Path PER> parameter is summation with a clamp of 10⁻¹ on the maximum value. The PERs are reported in units of 10^-11. A system with resolution less than 10^-11 MUST set unused digits to zero. An individual QNE adds its local PER value (up to a maximum of 10^{-2}) to the total Path PER value (up to a maximum of 10^{-1}), where the acceptability of the total Path PER value added across all QNEs is determined based on the QOSM being used. The maximum limit of 10^-2 on a QNE's local PER value and the maximum limit (clamp value) of 10^-1 on the accumulated end-to-end Path PER value are used to preserve the accuracy of the simple additive accumulation function specified and to avoid more complex accumulation functions. Furthermore, if these maximums are exceeded, then the path would likely not meet the QoS objectives. If the sum of the different elements' values exceeds 10^-1, the end-toend cumulative PER SHOULD be reported as indeterminate. A QNE that cannot accurately predict the PER of packets it is processing MUST raise the Not Supported N flag and either leave the value of Path PER as is, or add its best estimate of its lower bound. A raised Not Supported N flag indicates the value of Path PER is a lower bound of the real Path PER. The distinguished value 10⁻¹ is taken to mean indeterminate PER. A QNE that cannot accurately predict the PER of packets it is processing SHOULD set its local path PER parameter to this value. Because the composition function limits the composed sum to this value, receipt of this value at a network element or application indicates that the true path PER is not known. This MAY happen because one or more network elements could not supply a value or because the range of the composition calculation was exceeded.

Ash, et al. Experimental

[Page 42]

5.2.7. <Slack Term> Parameter

0	1	2	3
0 1 2 3 4 5 6	7 8 9 0 1 2 3 4	1567890123	845678901
+-+-+-+-+-+-+-+-++++++++-	+-+-+-+-+-+-+-+-	-+-+-+-+-+-+-+-+-+-	+-+-+-+++++++++++++++++++++++++++++++++
M E N r	7	r r r r	1
+-+-+-+-+-+-+-+-++	+-+-+-+-+-+-+-	-+	+-
Slack	Term (S) (32-b	oit unsigned intege	er)
+-	+ - + - + - + - + - + - + - + -	-+	+-+-+-++-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+

Slack term S MUST be nonnegative and is measured in microseconds [RFC2212]. The Slack term, S, is represented as a 32-bit unsigned integer. Its value can range from 0 to (2^32)-1 microseconds.

5.2.8. <Preemption Priority> and <Defending Priority> Parameters

The coding for the <Preemption Priority> and <Defending Priority> sub-parameters is as follows [RFC3181]:

0	1	2	3
0 1 2 3 4 5 6	7 8 9 0 1 2 3	4 5 6 7 8 9 0 1 2 3 4 5 6 7	78901
+-	+-+-+-+-+-+-+	-+	-+-+-+-+
M E N r	8	$ \mathbf{r} \mathbf{r} \mathbf{r} \mathbf{r} \qquad 1$	
+-	+-+-+-+-+-+-+	-+	-+-+-+-+
Preemption	n Priority	Defending Priorit	cy
+-	-+-+-+-+-+-+	-+	-+-+-+-+

- Preemption Priority: The priority of the new flow compared with the defending priority of previously admitted flows. Higher values represent higher priority.
- Defending Priority: Once a flow is admitted, the preemption priority becomes irrelevant. Instead, its defending priority is used to compare with the preemption priority of new flows.

As specified in [RFC3181], <Preemption Priority> and <Defending Priority> are 16-bit integer values, and both MUST be populated if the parameter is used.

Ash, et al. Experimental

[Page 43]

5.2.9. <Admission Priority> Parameter

The coding for the <Admission Priority> parameter is as follows:

0 1 2 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 $|\mathbf{M}|\mathbf{E}|\mathbf{N}|\mathbf{r}| \qquad 9 \qquad |\mathbf{r}|\mathbf{r}|\mathbf{r}| \qquad 1$ Y.2171 Adm Pri. Admis. Priority (Reserved)

Two fields are provided for the <Admission Priority> parameter and are populated according to the following rules.

<Y.2171 Admission Priority> values are globally significant on an end-to-end basis. High priority flows, normal priority flows, and best-effort priority flows can have access to resources depending on their admission priority value, as described in [Y.2171], as follows:

<Y.2171 Admission Priority>:

- 0 best-effort priority flow
- 1 normal priority flow
- 2 high priority flow

If the QNI signals <Y.2171 Admission Priority>, it populates both the <Y.2171 Admission Priority> and <Admission Priority> fields with the same value. Downstream QNEs MUST NOT change the value in the <Y.2171 Admission Priority> field so that end-to-end consistency is maintained and MUST treat the flow priority according to the value populated. A QNE in a local domain MAY reset a different value of <Admission Priority> in a Local QSPEC, but (as specified in Section 4.1) the Local QSPEC MUST be consistent with the Initiator QSPEC. That is, the local domain MUST specify an <Admission Priority> in the Local QSPEC that is functionally equivalent to the <Y.2171 Admission Priority> specified by the QNI in the Initiator QSPEC.

If the QNI signals admission priority according to [EMERGENCY-RSVP], it populates a locally significant value in the <Admission Priority> field and places all ones in the <Y.2171 Admission Priority> field. In this case, the functional significance of the <Admission Priority> value is specified by the local network administrator. Higher values indicate higher priority. Downstream QNEs and RSVP nodes MAY reset the <Admission Priority> value according to the local rules specified by the local network administrator, but MUST NOT reset the value of the <Y.2171 Admission Priority> field.

Ash, et al. Experimental

[Page 44]

A reservation without an <Y.2171 Admission Priority> parameter MUST be treated as a reservation with an <Y.2171 Admission Priority> = 1.

5.2.10. <RPH Priority> Parameter

The coding for the <RPH Priority> parameter is as follows:

0	1	2	3
0 1 2 3 4 5 6	7 8 9 0 1 2 3 4	5 6 7 8 9 0 1 2 3	4 5 6 7 8 9 0 1
+-+-+-+-+-+-	+-+-+-+-+-+-+-+-	+-	+-+-+-+-+-+-+-+
M E N r	10	r r r r	1
+-+-+-+-+-+-+-	+-+-+-+-+-+-+-+-+-+-+-+++	+-	+-
RPH	Namespace	RPH Priority	(Reserved)
+-+-+-+-+-+-	+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+++	+-	+-+-+-+-+-+-+-+

[RFC4412] defines a resource priority header (RPH) with parameters "RPH Namespace" and "RPH Priority", and if populated is applicable only to flows with high admission priority. A registry is created in [RFC4412] and extended in [EMERG-RSVP] for IANA to assign the RPH priority parameter. In the extended registry, "Namespace Numerical Values" are assigned by IANA to RPH Namespaces and "Priority Numerical Values" are assigned to the RPH Priority.

Note that the <Admission Priority> parameter MAY be used in combination with the <RPH Priority> parameter, which depends on the supported QOSM. Furthermore, if more than one RPH namespace is supported by a QOSM, then the QOSM MUST specify how the mapping between the priorities belonging to the different RPH namespaces are mapped to each other.

Note also that additional work is needed to communicate these flow priority values to bearer-level network elements [VERTICAL-INTERFACE].

For the 4 priority parameters, the following cases are permissible (procedures specified in references):

1	parameter:	<admission priority=""> [Y.2171]</admission>
2	parameters:	<admission priority="">, <rph priority=""> [RFC4412]</rph></admission>
2	parameters:	<pre><preemption priority="">, <defending priority=""> [RFC3181]</defending></preemption></pre>
3	parameters:	<preemption priority="">, <defending priority="">,</defending></preemption>
		<pre><admission priority=""> [3GPP-1, 3GPP-2, 3GPP-3]</admission></pre>
4	parameters:	<preemption priority="">, <defending priority="">,</defending></preemption>
		<pre><admission priority="">, <rph priority=""> [3GPP-1, 3GPP-2,</rph></admission></pre>
		3GPP-3]

Ash, et al. Experimental

[Page 45]

It is permissible to have <Admission Priority> without <RPH Priority>, but not permissible to have <RPH Priority> without <Admission Priority>. (Alternatively, <RPH Priority> is ignored in instances without <Admission Priority>.)

Functionality similar to enhanced Multi-Level Precedence and Preemption service (eMLPP; as defined in [3GPP-1, 3GPP-2]) specifies use of <Admission Priority> corresponding to the 'queuing allowed' part of eMLPP, as well as <Preemption/Defending Priority> corresponding to the 'preemption capable' and 'may be preempted' parts of eMLPP.

5.2.11. <Excess Treatment> Parameter

0	1	2	3
0 1 2 3 4 5 6	7890123	4 5 6 7 8 9 0 1 2 3 4	5678901
+-+-+-+-+-+-+	-+	-+	-+-+-+-+-+-+
M E N r	11	r r r r	1
+-	-+	-+	-+-+-+-+-+-+
Excess Trtmnt	Re-mark Val	Reserved	
+-	-+	-+	-+-+-+-+-+-+

Excess Treatment: Indicates how the QNE SHOULD process out-of-profile traffic, that is, traffic not covered by the <TMOD> parameter. The Excess Treatment Parameter is set by the QNI. Allowed values are as follows:

0: drop

- 1: shape
- 2: re-mark
- 3: no metering or policing is permitted

If no Excess Treatment Parameter is specified, the default is that there are no guarantees to excess traffic, i.e., a QNE can do whatever it finds suitable.

When excess treatment is set to 'drop', all marked traffic MUST be dropped by the QNE/RMF.

When excess treatment is set to 'shape', it is expected that the QoS Desired object carries a TMOD parameter, and excess traffic is shaped to this TMOD. The bucket size in the TMOD parameter for excess traffic specifies the queuing behavior, and when the shaping causes unbounded queue growth at the shaper, any traffic in excess of the TMOD for excess traffic SHOULD be dropped. If excess treatment is set to 'shape' and no TMOD parameter is given, the E flag is set for the parameter and the reservation fails. If

Ash, et al. Experimental

[Page 46]

excess treatment is set to 'shape' and two TMOD parameters are specified, then the QOSM specification dictates how excess traffic should be shaped in that case.

When excess treatment is set to 're-mark', the Excess Treatment Parameter MUST carry the re-mark value, and the re-mark values and procedures MUST be specified in the QOSM specification document. For example, packets may be re-marked to pertain to a particular QoS class (Diffserv Code Point (DSCP) value). In the latter case, re-marking relates to a Diffserv model where packets arrive marked as belonging to a certain QoS class / DSCP, and when they are identified as excess, they should then be re-marked to a different QoS Class (DSCP value) indicated in the 'Re-mark Value', as follows:

Re-mark Value (6 bits): indicates DSCP value [RFC2474] to re-mark packets to when identified as excess

If 'no metering or policing is permitted' is signaled, the QNE should accept the Excess Treatment Parameter set by the sender with special care so that excess traffic should not cause a problem. To request the Null Meter [RFC3290] is especially strong, and should be used with caution.

A NULL metering application [RFC2997] would not include the traffic profile, and conceptually it should be possible to support this with the QSPEC. A QSPEC without a traffic profile is not excluded by the current specification. However, note that the traffic profile is important even in those cases when the excess treatment is not specified, e.g., in negotiating bandwidth for the best-effort aggregate. However, a "NULL Service QOSM" would need to be specified where the desired QNE Behavior and the corresponding QSPEC format are described.

As an example behavior for a NULL metering, in the properly configured Diffserv router, the resources are shared between the aggregates by the scheduling disciplines. Thus, if the incoming rate increases, it will influence the state of a queue within that aggregate, while all the other aggregates will be provided sufficient bandwidth resources. NULL metering is useful for best-effort and signaling data, where there is no need to meter and police this data as it will be policed implicitly by the allocated bandwidth and, possibly, active queue management mechanism.

Ash, et al. Experimental

[Page 47]

5.2.12. <PHB Class> Parameter

The coding for the <PHB Class> parameter is as follows [RFC3140]:

0 2 1 3 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 |M|E|N|r| 12 |r|r|r|r| 1 PHB Field (Reserved)

The above encoding is consistent with [RFC3140], and the following four figures show four possible formats based on the value of the PHB Field.

Single PHB:

0 1 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 DSCP 0 0 0 0 0 0 0 0 0 0 0

Set of PHBs:

0 1 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 DSCP 000000010

PHBs not defined by standards action, i.e., experimental or local use PHBs as allowed by [RFC2474]. In this case, an arbitrary 12-bit PHB identification code, assigned by the IANA, is placed left-justified in the 16-bit field. Bit 15 is set to 1, and bit 14 is zero for a single PHB or 1 for a set of PHBs. Bits 12 and 13 are zero.

Single non-standard PHB (experimental or local):

0 1 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 PHB ID CODE 0 0 1

Ash, et al.

Experimental

[Page 48]

Set of non-standard PHBs (experimental or local):

1 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 PHB ID CODE 0 0 1 1

Bits 12 and 13 are reserved either for expansion of the PHB identification code, or for other use, at some point in the future.

In both cases, when a single PHBID is used to identify a set of PHBs (i.e., bit 14 is set to 1), that set of PHBs MUST constitute a PHB Scheduling Class (i.e., use of PHBs from the set MUST NOT cause intra-microflow traffic reordering when different PHBs from the set are applied to traffic in the same microflow). The set of AF1x PHBs [RFC2597] is an example of a PHB Scheduling Class. Sets of PHBs that do not constitute a PHB Scheduling Class can be identified by using more than one PHBID.

The registries needed to use RFC 3140 already exist; see [DSCP-REGISTRY] and [PHBID-CODES-REGISTRY]. Hence, no new registry needs to be created for this purpose.

5.2.13. <DSTE Class Type> Parameter

A description of the semantic of the parameter values can be found in [RFC4124]. The coding for the <DSTE Class Type> parameter is as follows:

0 2 1 3 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 $|\mathbf{M}|\mathbf{E}|\mathbf{N}|\mathbf{r}| \qquad 13 \qquad |\mathbf{r}|\mathbf{r}|\mathbf{r}|\mathbf{r}|$ 1 DSTE Cls. Type (Reserved)

DSTE Class Type: Indicates the DSTE class type. Values currently allowed are 0, 1, 2, 3, 4, 5, 6, and 7.

Ash, et al. Experimental

[Page 49]

5.2.14. <Y.1541 QoS Class> Parameter The coding for the <Y.1541 QoS Class> parameter [Y.1541] is as follows: 0 1 2 3 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 $|\mathbf{M}|\mathbf{E}|\mathbf{N}|\mathbf{r}| \qquad 14 \qquad |\mathbf{r}|\mathbf{r}|\mathbf{r}|\mathbf{r}|$ 1 Y.1541 QoS Cls. (Reserved) Y.1541 QoS Class: Indicates the Y.1541 QoS Class. Values currently allowed are 0, 1, 2, 3, 4, 5, 6, and 7. Class 0: Real-time, highly interactive applications, sensitive to jitter. Mean delay <= 100 ms, delay variation <= 50 ms, and loss ratio <= 10⁻³. Application examples include VoIP and video teleconference. Class 1: Real-time, interactive applications, sensitive to jitter. Mean delay <= 400 ms, delay variation <= 50 ms, and loss ratio <= 10^-3. Application examples include VoIP and video teleconference. Class 2: Highly interactive transaction data. Mean delay <= 100 ms, delay variation is unspecified, loss ratio <= 10^-3. Application examples include signaling. Class 3: Interactive transaction data. Mean delay <= 400 ms, delay variation is unspecified, loss ratio <= 10^-3. Application examples include signaling. Class 4: Low Loss Only applications. Mean delay <= 1 s, delay variation is unspecified, loss ratio <= 10^-3. Application examples include short transactions, bulk data, and video streaming. Class 5: Unspecified applications with unspecified mean delay, delay variation, and loss ratio. Application examples include traditional applications of default IP networks. Ash, et al. Experimental [Page 50]

Class 6: Applications that are highly sensitive to loss. Mean delay <= 100 ms, delay variation <= 50 ms, and loss ratio <= 10^{-5} . Application examples include television transport, high-capacity TCP transfers, and Time-Division Multiplexing (TDM) circuit emulation.

Class 7: Applications that are highly sensitive to loss. Mean delay <= 400 ms, delay variation <= 50 ms, and loss ratio <= 10⁻⁵. Application examples include television transport, high-capacity TCP transfers, and TDM circuit emulation.

6. Security Considerations

QSPEC security is directly tied to QoS NSLP security, and the QoS NSLP document [RFC5974] has a very detailed security discussion in Section 7. All the considerations detailed in Section 7 of [RFC5974] apply to QSPEC.

The priority parameter raises possibilities for theft-of-service attacks because users could claim an emergency priority for their flows without real need, thereby effectively preventing serious emergency calls to get through. Several options exist for countering such attacks, for example:

- only some user groups (e.g., the police) are authorized to set the emergency priority bit
- any user is authorized to employ the emergency priority bit for particular destination addresses (e.g., police)
- 7. IANA Considerations

This section defines the registries and initial codepoint assignments for the QSPEC template, in accordance with BCP 26, RFC 5226 [RFC5226]. It also defines the procedural requirements to be followed by IANA in allocating new codepoints.

This specification creates the following registries with the structures as defined below:

Object Types (12 bits): The following values are allocated as specified in Section 5: 0: QoS Desired 1: QoS Available 2: OoS Reserved 3: Minimum QoS

Ash, et al.

Experimental

[Page 51]

Further values are as follows: 4-63: Unassigned 64-67: Private/Experimental Use 68-4095: Reserved (Note: 'Reserved' just means 'do not give these out'.) The registration procedure is Specification Required. QSPEC Version (4 bits): The following value is allocated by this specification: 0: Version 0 QSPEC Further values are as follows: 1-15: Unassigned The registration procedure is Specification Required. (A specification is required to depreciate, delete, or modify QSPEC versions.) QSPEC Type (5 bits): The following values are allocated by this specification: 0: Default 1: Y.1541-QOSM [RFC5976] 2: RMD-QOSM [RFC5977] Further values are as follows: 3-12: Unassigned 13-16: Local/Experimental Use 17-31: Reserved The registration procedure is Specification Required. QSPEC Procedure (8 bits): The QSPEC Procedure object consists of the Message Sequence parameter

(4 bits) and the Object Combination parameter (4 bits), as discussed in Section 4.3. Message Sequences 0 (Two-Way Transactions), 1 (Three-Way Transactions), and 2 (Resource Queries) are explained in Sections 4.3.1, 4.3.2, and 4.3.3, respectively. Tables 1, 2, and 3 in Section 4.3 assign the Object Combination Number to Message Sequences 0, 1, and 2, respectively. The values assigned by this specification for the Message Sequence parameter and the Object Combination parameter are summarized here:

Ash, et al.

Experimental

[Page 52]

RFC	6075	
RFC	5915	

	!		OBJECTS INCLUDED IN RESERVE MESSAGE	OBJECTS INCLUDED IN RESPONSE MESSAGE
0	0	N/A	QOS Desired	QoS Reserved
0	1	N/A N/A	 QoS Desired QoS Available	 QoS Reserved QoS Available
0	2 	N/A N/A N/A	QoS Desired QoS Available Minimum QoS	QoS Reserved QoS Available
1	0	QoS Desired	QoS Desired	QoS Reserved
1	1	QoS Desired (Minimum QoS)	QoS Desired QoS Available (Minimum QoS)	QoS Reserved QoS Available
1	2	QoS Desired QoS Available	 QoS Desired QoS Available	QoS Reserved
2	0	QoS Available	N/A	QoS Available

Further values of the Message Sequence parameter (4 bits) are as follows:

3-15: Unassigned

Further values of the Object Combination parameter (4 bits) are as follows:

Message | Object Sequence | Combination -----0 | 3-15: Unassigned 1 2 3-15: Unassigned | 1-15: Unassigned 3-15 0-15: Unassigned

The registration procedure is Specification Required. (A specification is required to depreciate, delete, or modify QSPEC Procedures.)

QoS Model Error Code (8 bits): QoS Model Error Codes may be defined for NSLP error class 6 (QoS Model Error), as described in Section 6.4 of [RFC5974]. Values are as follows: 0-63: Unassigned 64-67: Private/Experimental Use

Ash, et al.

Experimental

[Page 53]

```
RFC 5975
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```
68-255: Reserved
   The registration procedure is Specification Required. (A
   specification is required to depreciate, delete, or modify QoS Model
   Error Codes.)
   Parameter ID (12 bits):
   The following values are allocated by this specification:
   1-14: assigned as specified in Section 5.2:
     1: <TMOD-1>
      2: <TMOD-2>
     3: <Path Latency>
      4: <Path Jitter>
      5: <Path PLR>
      6: <Path PER>
      7: <Slack Term>
      8: <Preemption Priority> and <Defending Priority>
      9: <Admission Priority>
     10: <RPH Priority>
     11: <Excess Treatment>
     12: <PHB Class>
     13: <DSTE Class Type>
     14: <Y.1541 QoS Class>
   Further values are as follows:
      15-255: Unassigned
      256-259: Private/Experimental Use
      260-4095: Reserved
   The registration procedure is Specification Required. (A
   specification is required to depreciate, delete, or modify Parameter
   IDs.)
   Y.2171 Admission Priority Parameter (8 bits):
   The following values are allocated by this specification:
   0-2: assigned as specified in Section 5.2.9:
      0: best-effort priority flow
      1: normal priority flow
      2: high priority flow
   Further values are as follows:
      3-63: Unassigned
      64-255: Reserved
   The registration procedure is Specification Required.
   RPH Namespace Parameter (16 bits):
  Note that [RFC4412] creates a registry for RPH Namespace and Priority
   values already (see Section 12.6 of [RFC4412]), and an extension to
   this registry is created in [EMERG-RSVP], which will also be used for
   the QSPEC RPH parameter. In the extended registry, "Namespace
   Numerical Values" are assigned by IANA to RPH Namespaces, and
            Experimental
Ash, et al.
                                                               [Page 54]
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"Priority Numerical Values" are assigned to the RPH Priority. There are no additional IANA requirements made by this specification for the RPH Namespace Parameter. Excess Treatment Parameter (8 bits): The following values are allocated by this specification: 0-3: assigned as specified in Section 5.2.11: 0: drop 1: shape 2: re-mark 3: no metering or policing is permitted Further values are as follows: 4-63: Unassigned 64-255: Reserved The registration procedure is Specification Required. Y.1541 QoS Class Parameter (8 bits): The following values are allocated by this specification: 0-7: assigned as specified in Section 5.2.14: 0: Y.1541 QoS Class 0 1: Y.1541 QoS Class 1 2: Y.1541 QoS Class 2 3: Y.1541 QoS Class 3 4: Y.1541 QoS Class 4 5: Y.1541 QoS Class 5 6: Y.1541 QoS Class 6 7: Y.1541 QoS Class 7 Further values are as follows: 8-63: Unassigned 64-255: Reserved The registration procedure is Specification Required.

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Ash, et al.

Experimental

[Page 55]

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Ash, et al.

Experimental

[Page 56]

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Ash, et al.

Experimental

[Page 61]

Appendix A. Mapping of QoS Desired, QoS Available, and QoS Reserved of NSIS onto AdSpec, TSpec, and RSpec of RSVP IntServ

The union of QoS Desired, QoS Available, and QoS Reserved can provide all functionality of the objects specified in RSVP IntServ; however, it is difficult to provide an exact mapping.

In RSVP, the Sender TSpec specifies the traffic an application is going to send (e.g., TMOD). The AdSpec can collect path characteristics (e.g., delay). Both are issued by the sender. The receiver sends the FlowSpec that includes a Receiver TSpec describing the resources reserved using the same parameters as the Sender TSpec, as well as an RSpec that provides additional IntServ QoS Model specific parameters, e.g., Rate and Slack.

The RSVP TSpec, AdSpec, and RSpec are tailored to the receiverinitiated signaling employed by RSVP and the IntServ QoS Model. For example, to the knowledge of the authors, it is not possible for the sender to specify a desired maximum delay except implicitly and mutably by seeding the AdSpec accordingly. Likewise, the RSpec is only meaningfully sent in the receiver-issued RSVP RESERVE message. For this reason, our discussion at this point leads us to a slightly different mapping of necessary functionality to objects, which should result in more flexible signaling models.

Appendix B. Example of TMOD Parameter Encoding

In an example VoIP application that uses RTP [RFC3550] and the G.711 Codec [G.711], the TMOD-1 parameter could be set as follows:

In the simplest case, the Minimum Policed Unit m is the sum of the IP, UDP, and RTP headers + payload. The IP header in the IPv4 case has a size of 20 octets (40 octets if IPv6 is used). The UDP header has a size of 8 octets, and RTP uses a 12-octet header. The G.711 Codec specifies a bandwidth of 64 kbit/s (8000 octets/s). Assuming RTP transmits voice datagrams every 20 ms, the payload for one datagram is 8000 octets/s * 0.02 s = 160 octets.

IPv4 + UDP + RTP + payload: m = 20 + 8 + 12 + 160 octets = 200 octets IPv6 + UDP + RTP + payload: m = 40 + 8 + 12 + 160 octets = 220 octets

The Rate r specifies the amount of octets per second. 50 datagrams are sent per second.

IPv4: $r = 50 \ 1/s * m = 10,000 \text{ octets/s}$ IPv6: r = 50 1/s * m = 11,000 octets/s

Ash, et al. Experimental

[Page 62]

The bucket size b specifies the maximum burst. In this example, a burst of 10 packets is used.

IPv4: b = 10 * m = 2000 octets IPv6: b = 10 * m = 2200 octets

A number of extra headers (e.g., for encapsulation) may be included in the datagram. A non-exhaustive list is given below. For additional headers, m, r, and b have to be set accordingly.

Protocol Header Size

	+
GRE [RFC1701]	8 octets
GREIP4 [RFC1702]	4-8 octets
IP4INIP4 [RFC2003]	20 octets
MINENC [RFC2004]	8-12 octets
IP6GEN [RFC2473]	40 octets
IP6INIP4 [RFC4213]	20 octets
IPsec [RFC4301, RFC4303]	variable
	+

Experimental

[Page 63]

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Ash, et al.

Experimental

[Page 64]